

# 即時控制系統設計

## Design of Real-Time Control Systems

### Lecture 32

#### Introduction to Networked Control Systems

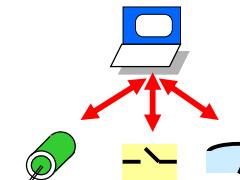
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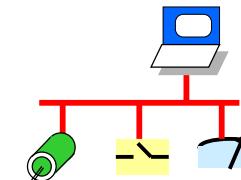
Feb10 – Jun10

### ▪ Real-Time Control Systems

- Controlled by one Computer Processor
  - Centralized control systems
  - Real-time operating systems
- Controlled by one Communication Medium
  - Distributed control systems
  - Real-time communications



Centralized Control System



Distributed Control System

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### ▪ Networks:

- Enable remote data transfers & data exchange among users
- Reduce the complexity in wiring connections and the costs of media
- Provide ease in maintenance

#### • Data Networks:

- > Slotted ALOHA & ARPANET around 1960-70
- > Ethernet around 1980

#### • Control Networks:

> CAN (Controller Area Network) in 1983:

» By Robert Bosch, Germany, for car industries

> PROFIBUS (PROcess FieId BUS) in 1987:

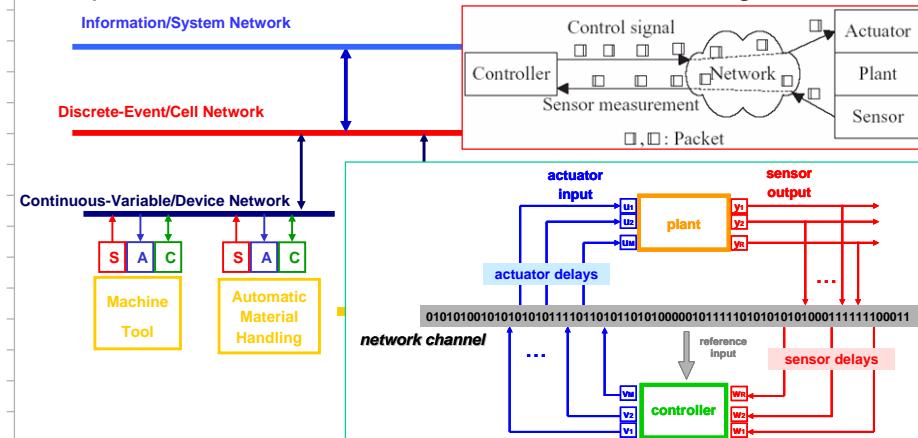
» By six German companies & five German institutes

> DeviceNet in 1994 (?)

» By Allen-Bradley/Rockwell Automation, for manufacturing-related industries

### ▪ Networked Control Systems:

- Control systems with physically distributed processing power and network communication of control signals



## Motivations:

- The overall NCS performance is always affected by **network delays**
- **Delays** are widely known to **degrade** the performance of a control system
- Existing **constant time-delay** control methodologies may **not** be directly **suitable** for controlling a system over the network since **network delays** are usually **time-varying**, especially in the Internet
- Therefore, to handle **network delays** in a closed-loop control system over a network, an advanced methodology is required

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## Overview of NCS Research Issues

### Research Overview:

- **NCS Configuration**
  - » Direct structure
  - » Hierarchical structure
- **Delays in-the-loop**
- **Delay Characteristics**
  - » Cyclic service network
  - » Random access network
- **Effect of Delays in-the-loop**
  - > Performance degradation
  - > Destabilization

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## Overview of NCS Research Issues

### NCS Configuration:

#### Direct structure:

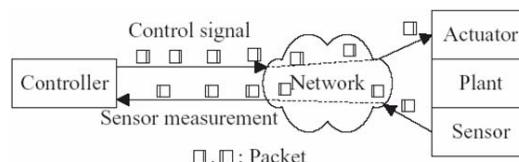


Fig. 1. NCS in the direct structure.

#### Hierarchical structure:

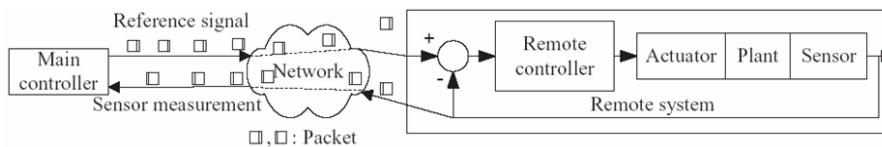


Fig. 2. NCS in the hierarchical structure.

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## Overview of NCS Research Issues

### NCS Configuration:

#### Direct structure:

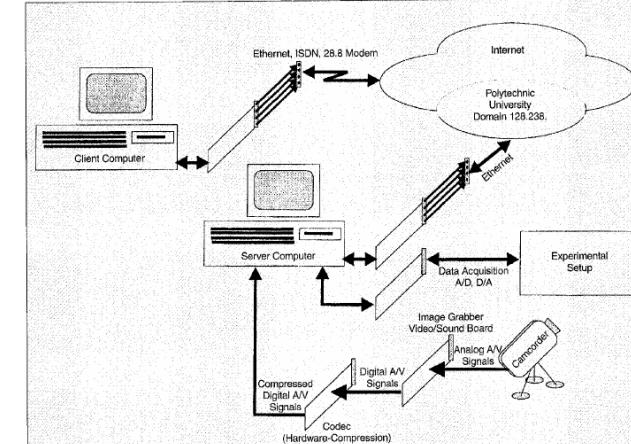


Fig. 9. I-abc network topology/experiment configuration.

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## ▪ NCS Configuration:

- Direct structure:

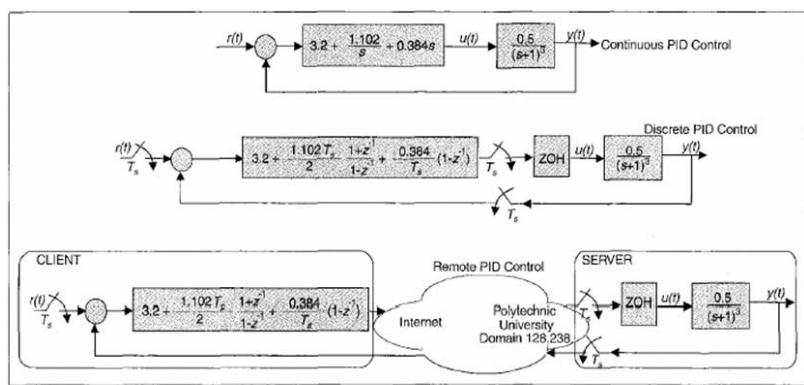


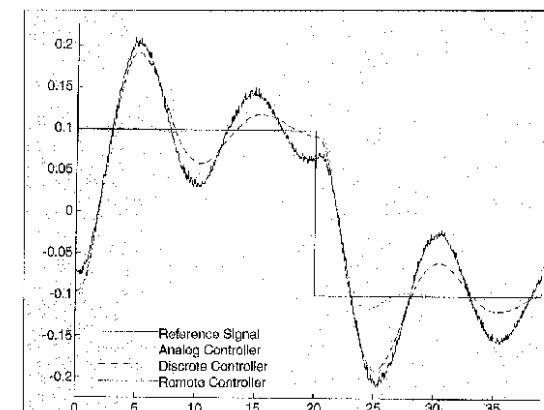
Fig. 10. Continuous, discrete and remote PID control concepts.

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## ▪ NCS Configuration:

- Direct structure:

Fig. 13. System response with continuous, discrete and remote ( $T_s = 2000 \text{ ns}$ ) control.

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## ▪ NCS Configuration:

- Hierarchical structure:

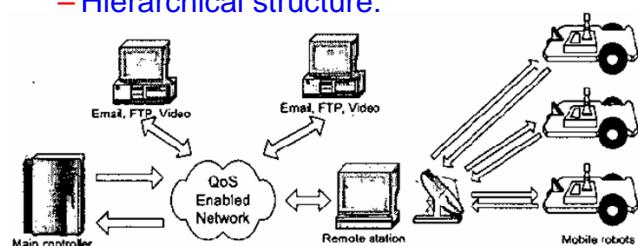


Fig. 1. Distributed mobile robot system configuration

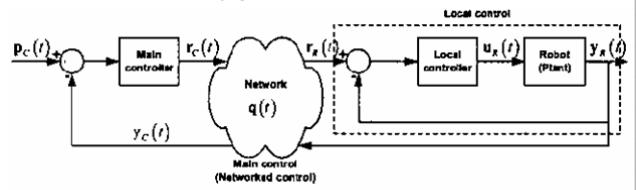


Fig. 2. Block diagram of main and local controllers

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## ▪ Delays in-the-loop:

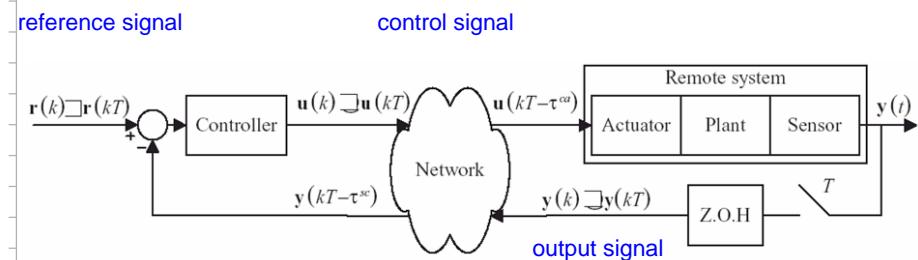


Fig. 3. General NCS configuration and network delays for NCS formulations.

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## ▪ Delays in-the-loop:

$\tau^{sc}$ : sensor-to-controller delay

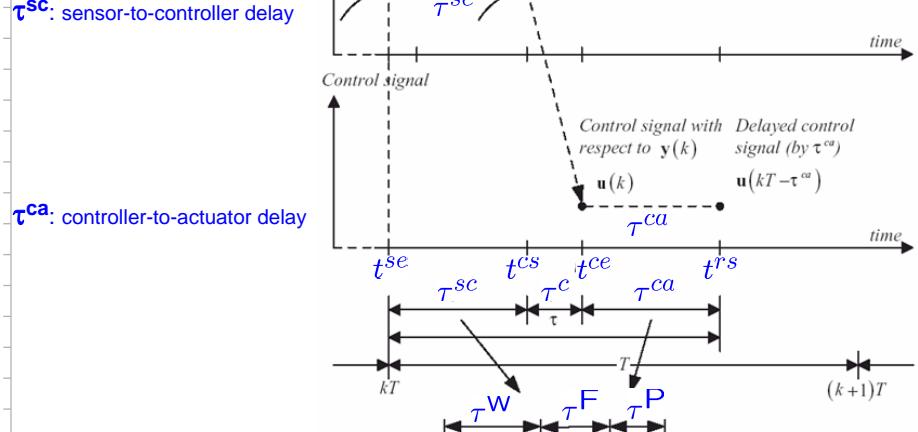
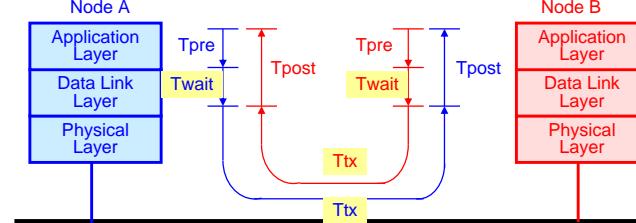


Fig. 4. Timing diagram of network delay propagations.

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## ▪ Delays in-the-loop:



- Total end-to-end delay is the sum of
  - Pre-processing time: microprocessor
  - Waiting time: network protocol - MAC
  - Transmission time: data rate & length
  - Post-processing time: microprocessor

Depend on network protocol and loading

## ▪ Delay Characteristics:

### • Cyclic Service Networks:

- IEEE 802.4, SAE token bus, PROFIBUS, IEEE 802.5, SAE token ring, MIL-STD-1553B, FIP
- Control and sensory signals are transmitted in a cyclic order with deterministic behaviors
- Delays are periodic & can be simply modeled as a periodic function such as  $\tau_k^{sc} = \tau_{k+N}^{sc}$  and  $\tau_k^{ca} = \tau_{k+M}^{ca}$ ,  $N, M$  are constants
- In practice, NCS may experience small variations on periodic delays due to several reasons such as discrepancies in clock generators

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## ▪ Delay Characteristics:

### • Random Access Networks:

- Ethernet, CAN (?)
- Significant parts of random network delays are waiting time delay due to queuing and frame collision on the networks

### – Sources of randomness:

- > the queuing time delay at a switch or a router
- > The propagation time delays from different network paths

### – Delay models:

- > Constant delay
- > Poisson process such as Markov chain
- > Fluid flow model, ARMA model etc.

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## Overview of NCS Research Issues

### ▪ Delay Models:

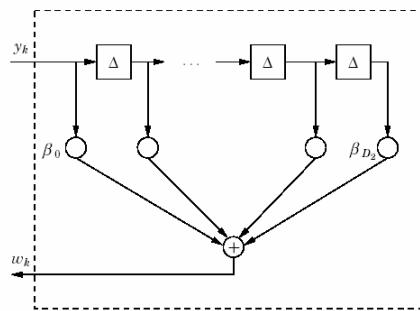


Figure 2.5 In Luck and Ray (1990) buffers are introduced after the varying communication delays to make the system time-invariant. The buffers must be longer than the worst case communication delay.

Figure 2.6 The network model used in Krtolica *et al.* (1994). The sampled signal,  $y_k$ , is delayed a number of samples due to communication delay. The controller reads the signal  $w_k$ . Only one of the  $\beta_i$  coefficients is one, the others are zero.

Luck & Ray 90, Krtolica *et al.* 94, Nilsson 98

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## Overview of NCS Research Issues

### ▪ Delay Models:

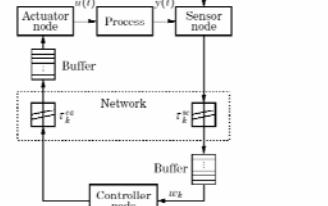


Figure 2.5 In Luck and Ray (1990) buffers are introduced after the varying communication delays to make the system time-invariant. The buffers must be longer than the worst case communication delay.

Figure 2.7 Block diagram of the transmission from the sensor node to the controller node in Chan and Ozguner (1995). The sampled signal  $y_k$  is delayed during the transmission to the controller node. The controller reads the sensor value  $w_k$  from a register in the controller node. A simple form of timestamping is done by appending every message with the size of the queue when the message was sent.

Luck & Ray 90, Krtolica *et al.* 94, Chan & Ozguner 95, Nilsson 98

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## Overview of NCS Research Issues

### ▪ Delay Models:

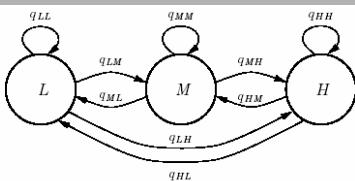


Figure 3.1 An example of a Markov chain modeling the state in a communication network. L is the state for low network load, M is the state for medium network load, and H is the state for high network load. The arrows show possible transitions in the system.

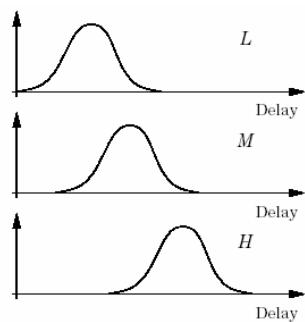


Figure 3.2 The delay distributions corresponding to the states of the Markov chain in Figure 3.1. L is the state for low network load, M is the state for medium network load, and H is the state for high network load.

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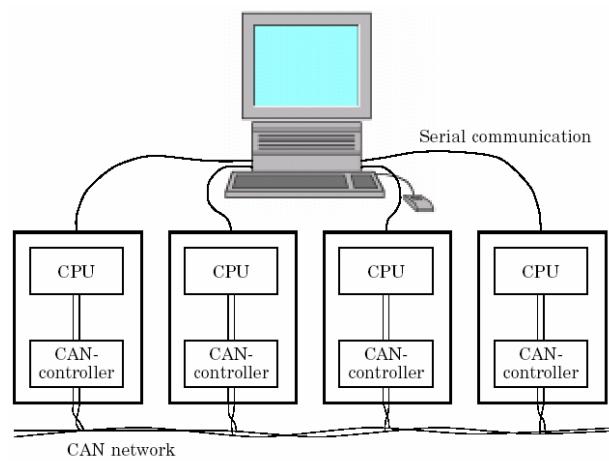
$$f_{\tau}(\tau_k^{ca}) = \delta(\tau_k^{ca} - a) \cdot (1 - p_{ca}) + \delta(\tau_k^{ca} - b) \cdot p_{ca}$$

$$f_{\tau}(\tau_k^{sc}) = \begin{cases} \delta(\tau_k^{sc} - a) \cdot (1 - p_{sc}), & \tau_k^{sc} = a, \\ p_{sc}/(b - a^+), & \tau_k^{sc} \in (a, b], a < b, \\ 0, & \tau_k^{sc} \notin (a, b], \end{cases}$$

Tipsuwan & Chow 03, Krtolica *et al.* 94, Nilsson 98

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## ▪ Experimental Data of Network Delays:



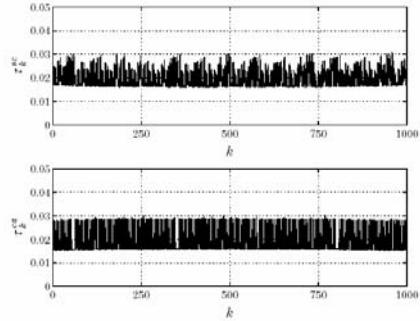
**Figure 4.1** The CAN experimental system with the four one-card computers. The nodes are connected with a twisted pair CAN. Serial communications to the nodes are used for data collection.

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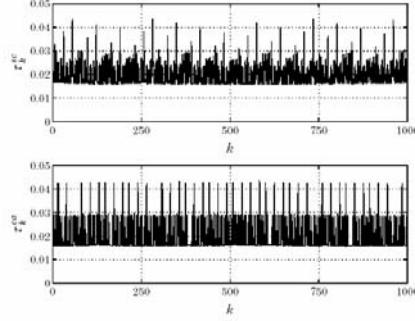
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## ▪ Experimental Data of Network Delays:

- Add two additional low-load nodes with different priorities



**Figure 4.7** Load: Two senders with periods 80 ms and 100 ms and lower pri. Message length: 8 bytes.



**Figure 4.8** Load: Two load processes, one with period 80 ms and higher priority, and one with period 100 ms and lower priority. Message length: 8 bytes.

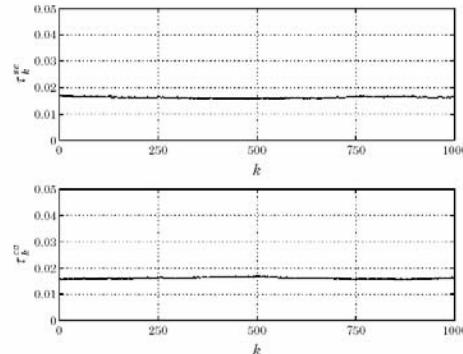
$$U = \frac{0.0128}{0.05} + \frac{0.0128}{0.05} + \frac{0.0128}{0.08} + \frac{0.0128}{0.1} = 0.80.$$

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## ▪ Experimental Data of Network Delays:

- Low load network traffic



**Figure 4.3** Measured delays for sensor to controller,  $\tau_k^{sc}$ , and from controller to actuator,  $\tau_k^{ac}$ . The experiment was the only load on the bus. The messages were 8 bytes long. The delay is almost constant.

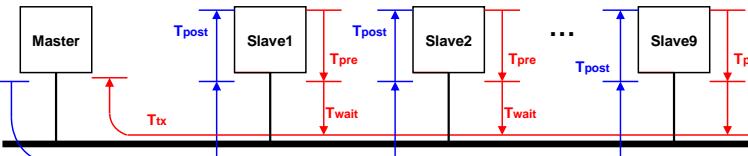
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## ▪ Experimental Data of Network Delays:

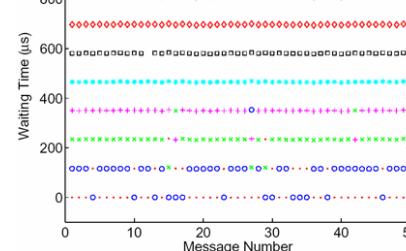
## ▪ Experimental Data of Network Delays:

- Nine-node network



**Figure 4.9** Nine-node network

**Figure 4.10** Waiting time vs. message number



- Period: 50 ms
- Data rate: 10 kbps
- Data size: 8 bytes/128 bits
- Tx time: 12.8 ms
- CPU time: 3.4 ms

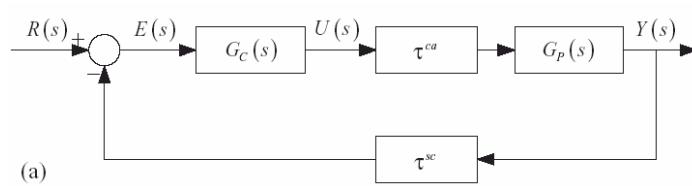
$$U = \frac{0.0128}{0.05} + \frac{0.0128}{0.05} = 0.51.$$

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▪ Effects of Delays in-the-loop:

- Performance degradation:



$$G_C(s) = \frac{\beta K_P (s + (K_I/K_P))}{s}, \quad G_P(s) = \frac{2029.826}{(s + 26.29)(s + 2.296)},$$

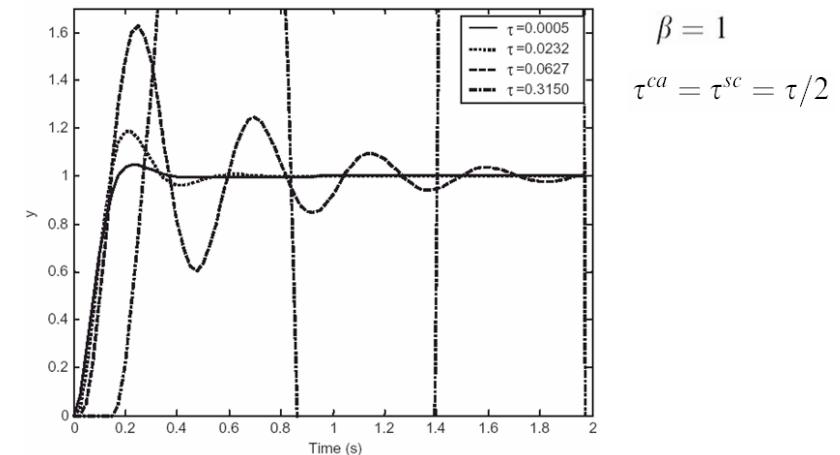
$$K_P = 0.1701, \quad K_I = 0.378,$$

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▪ Effects of Delays in-the-loop:

- Performance degradation:

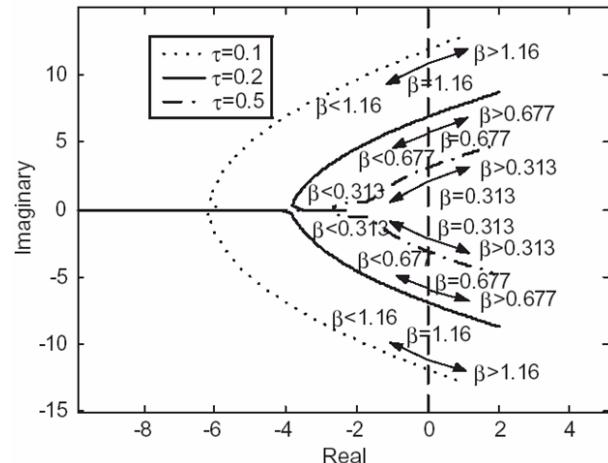


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▪ Effects of Delays in-the-loop:

- Destabilization: Root locus approach

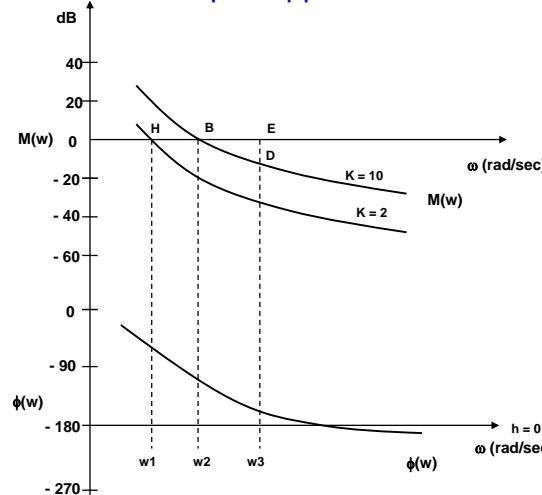


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▪ Effects of Delays in-the-loop:

- Destabilization: Bode plot approach

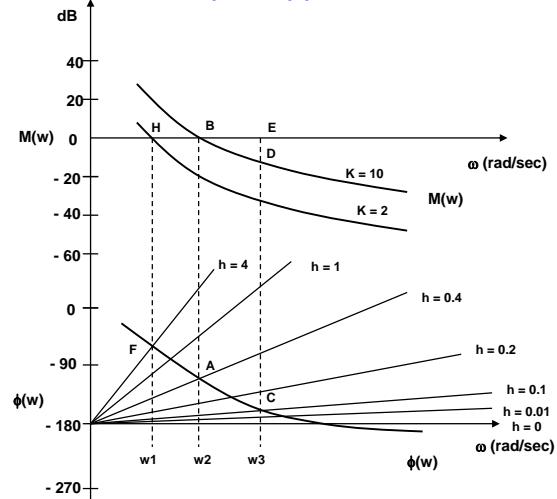


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- Effects of Delays in-the-loop:

- Destabilization: Bode plot approach



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- Networked Control Methodology:

- Assumptions
  - 1. Augmented Deterministic Discrete-Time Model Methodology
  - 2. Queuing Methodology
  - 3. Optimal Stochastic Control Methodology
  - 4. Perturbation Methodology
  - 5. Sampling Time Scheduling Methodology
  - 6. Robust Control Methodology
  - 7. Fuzzy Logic Modulation Methodology
  - 8. Event-Based Methodology
  - 9. End-User Control Adaptation Methodology

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