

Logic Synthesis and Verification

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Boolean Algebra

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Boolean Algebra

□ Reading

F. M. Brown. *Boolean Reasoning: The Logic of Boolean Equations*. Dover, 2003.
(Chapters 1-3)

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Boolean Algebra

□ Outline

- Definitions
- Examples
- Properties
- Boolean formulae and Boolean functions

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Boolean Algebra

□ A Boolean algebra is an algebraic structure $(\mathbf{B}, +, \cdot, \underline{0}, \underline{1})$

- $\mathbf{B}$ is a set, called the *carrier*
- $+$ and $\cdot$ are binary operations defined on $\mathbf{B}$
- $\underline{0}$ and $\underline{1}$ are distinct members of $\mathbf{B}$

that satisfies the following postulates (axioms):

1. *Commutative laws*
2. *Distributive laws*
3. *Identities*
4. *Complements*

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Postulates of Boolean Algebra

$(\mathbf{B}, +, \cdot, \underline{0}, \underline{1})$

1. $\mathbf{B}$ is **closed** under $+$ and $\cdot$
 $\forall a, b \in \mathbf{B}, a + b \in \mathbf{B}$ and $a \cdot b \in \mathbf{B}$
2. **Commutative laws**: $\forall a, b \in \mathbf{B}$
 $a + b = b + a$
 $a \cdot b = b \cdot a$
3. **Distributive laws**: $\forall a, b \in \mathbf{B}$
 $a + (b \cdot c) = (a + b) \cdot (a + c)$
 $a \cdot (b + c) = a \cdot b + a \cdot c$
4. **Identities**: $\forall a \in \mathbf{B}$
 $\underline{0} + a = a$
 $\underline{1} \cdot a = a$
5. **Complements**: $\forall a \in \mathbf{B}, \exists a' \in \mathbf{B}$ s.t.
 $a + a' = \underline{1}$
 $a \cdot a' = \underline{0}$
Verify that a' is unique in $(\mathbf{B}, +, \cdot, \underline{0}, \underline{1})$.

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Instances of Boolean Algebra

- Switching algebra (two-element Boolean algebra)
- The algebra of classes (subsets of a set)
- Arithmetic Boolean algebra
- The algebra of propositional functions

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Instance 1: Switching Algebra

□ A switching algebra is a two-element Boolean Algebra $(\{0,1\}, +, \cdot, 0, 1)$ consisting of:

- the set $\mathbf{B} = \{0, 1\}$
- two binary operations AND($\cdot$) and OR($+$)
- one unary operation NOT($'$)

where

OR	0	1
0	0	1
1	1	1

AND	0	1
0	0	0
1	0	1

NOT	-
0	1
1	0

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Switching Algebra

- Just one of many other Boolean algebras
 - (Ex: verify that the algebra satisfies all the postulates.)

- An exclusive property (not hold for all Boolean algebras) for two-element Boolean algebra:

$$x + y = 1 \text{ iff } x=1 \text{ or } y=1$$

$$x \cdot y = 0 \text{ iff } x=0 \text{ or } y=0$$

OR	0	1
0	0	1
1	1	1

AND	0	1
0	0	0
1	0	1

NOT	-
0	1
1	0

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Instance 2: Algebra of Classes

- Subsets of a set

$$\mathbf{B} \leftrightarrow 2^S$$

$$+ \leftrightarrow \cup$$

$$\cdot \leftrightarrow \cap$$

$$\underline{0} \leftrightarrow \phi$$

$$\underline{1} \leftrightarrow S$$

- S is a universal set ($S \neq \phi$). Each subset of S is called a *class* of S .

- If $S = \{a, b\}$, then $\mathbf{B} = \{\phi, \{a\}, \{b\}, \{a, b\}\}$

- $\mathbf{B}$ ($= 2^S$) is **closed** under $\cup$ and $\cap$

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Algebra of Classes

- Commutative laws:** $\forall S_1, S_2 \in 2^S$

$$S_1 \cup S_2 = S_2 \cup S_1$$

$$S_1 \cap S_2 = S_2 \cap S_1$$

- Distributive laws:** $\forall S_1, S_2, S_3 \in 2^S$

$$S_1 \cup (S_2 \cap S_3) = (S_1 \cup S_2) \cap (S_1 \cup S_3)$$

$$S_1 \cap (S_2 \cup S_3) = (S_1 \cap S_2) \cup (S_1 \cap S_3)$$

- Identities:** $\forall S_1 \in 2^S$

$$S_1 \cup \phi = S_1$$

$$S_1 \cap S = S_1$$

- Complements:** $\forall S_1 \in 2^S, \exists S_1' \in 2^S, S_1' = S \setminus S_1$ s.t.

$$S_1 \cup S_1' = S$$

$$S_1 \cap S_1' = \phi$$

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Algebra of Classes

- Stone Representation Theorem:**

Every finite Boolean algebra is isomorphic to the Boolean algebra of subsets of some finite set S

Therefore, for all finite Boolean algebra, $|\mathbf{B}|$ can only be 2^k for some $k \geq 1$.

- The theorem proves that finite class algebras are not specialized (i.e. no exclusive properties, e.g. $x + y = 1$ iff $x=1$ or $y=1$ in two-element Boolean algebra)

- Can reason in terms of specific and easily “visualizable” concepts (union, intersection, empty set, universal set) rather than abstract operations ($+$, $\cdot$, $\underline{0}$, $\underline{1}$)

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Instance 3: Arithmetic Boolean Algebra

- $(D_n, lcm, gcd, 1, n)$
 - n : product of distinct prime numbers
 - D_n : set of all divisors of n
 - lcm : least common multiple
 - gcd : greatest common divisor
 - 1: integer 1 (not the Boolean 1-element)
- $n = 30 = 2 \times 3 \times 5$
- $D_n = \{1, 2, 3, 5, 6, 10, 15, 30\}$
- If we look at D_n as $\{\emptyset, \{2\}, \{3\}, \{5\}, \{2, 3\}, \{2, 5\}, \{3, 5\}, \{2, 3, 5\}\}$, it is easy to see that arithmetic Boolean algebra is isomorphic to the algebra of classes.
 - See Stone Representation Theorem

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Instance 4: Algebra of Propositional Functions

- $(P, \vee, \wedge, \square, \blacksquare)$
 - P : the set of propositional functions of n given variables
 - $\vee$: disjunction symbol (OR)
 - $\wedge$: conjunction symbol (AND)
 - $\square$: formula that is always false (contradiction)
 - $\blacksquare$: formula that is always true (tautology)

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Lessons from Abstraction

- Abstract mathematical objects in terms of simple rules
- A systematic way of characterizing various seemingly unrelated mathematical objects
- Abstraction trims off immaterial details and simplifies problem formulation

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Properties of Boolean Algebras

- For arbitrary elements a , b , and c in Boolean algebra
- | | |
|---|---|
| 1. Associativity
$a + (b + c) = (a + b) + c$
$a \cdot (b \cdot c) = (a \cdot b) \cdot c$ | 5. Involution
$(a')' = a$ |
| 2. Idempotence
$a + a = a$
$a \cdot a = a$ | 6. De Morgan's Laws
$(a + b)' = a' \cdot b'$
$(a \cdot b)' = a' + b'$ |
| 3.
$a + \underline{1} = \underline{1}$
$a \cdot \underline{0} = \underline{0}$ | 7.
$a + a' \cdot b = a + b$
$a \cdot (a' + b) = a \cdot b$ |
| 4. Absorption
$a + (a \cdot b) = a$
$a \cdot (a + b) = a$ | 8. Consensus
$a \cdot b + a' \cdot c + b \cdot c =$
$a \cdot b + a' \cdot c$
$(a + b) \cdot (a' + c) \cdot (b + c) =$
$(a + b) \cdot (a' + c)$ |

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Principle of Duality

□ Every identity on Boolean algebra is transformed into another identity if the following is interchanged

- the operations $+$ and $\cdot$,
- the elements $\underline{0}$ and $\underline{1}$

□ Example:

- $a + \underline{1} = \underline{1}$
- $a \cdot \underline{0} = \underline{0}$

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Postulates for Boolean Algebra (Revisited)

Duality in $(\mathbf{B}, +, \cdot, \underline{0}, \underline{1})$

1. $\mathbf{B}$ is **closed** under $+$ and $\cdot$.
 $\forall a, b \in \mathbf{B}, a + b \in \mathbf{B}$ and $a \cdot b \in \mathbf{B}$
2. **Commutative Laws:** $\forall a, b \in \mathbf{B}$
 $a + b = b + a$
 $a \cdot b = b \cdot a$
3. **Distributive laws:** $\forall a, b \in \mathbf{B}$
 $a + (b \cdot c) = (a + b) \cdot (a + c)$
 $a \cdot (b + c) = a \cdot b + a \cdot c$
4. **Identities:** $\forall a \in \mathbf{B}$
 $\underline{0} + a = a$
 $\underline{1} \cdot a = a$
5. **Complements:** $\forall a \in \mathbf{B}, \exists a' \in \mathbf{B}$ s.t.
 $a + a' = \underline{1}$
 $a \cdot a' = \underline{0}$

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Two Propositions

1. Let a and b be members of a Boolean algebra. Then
 $a = \underline{0}$ and $b = \underline{0}$ iff $a + b = \underline{0}$
 $a = \underline{1}$ and $b = \underline{1}$ iff $ab = \underline{1}$

c.f. The following two propositions are only true for two-element Boolean algebra (not other Boolean algebra)

$$x + y = 1 \text{ iff } x = 1 \text{ or } y = 1$$
$$xy = 0 \text{ iff } x = 0 \text{ or } y = 0$$

Why?

2. Let a and b be members of a Boolean algebra. Then
 $a = b$ iff $a'b + ab' = \underline{0}$

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Boolean Formulas and Boolean Functions

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Boolean Formulas and Boolean Functions

□ Outline:

- Definition of Boolean formulas
- Definition of Boolean functions
- Boole's expansion theorem
- The minterm canonical form

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n -variable Boolean Formulas

- Given a Boolean algebra $\mathbf{B}$ and n symbols $x_1, \dots, x_n$, the set of all Boolean formulas on the n symbols is defined by:
 1. The elements of $\mathbf{B}$ are Boolean formulas.
 2. The variable symbols $x_1, \dots, x_n$ are Boolean formulas.
 3. If g and h are Boolean formulas, then so are
 - $(g) + (h)$
 - $(g) \cdot (h)$
 - $(g)'$
 4. A string is a Boolean formula if and only if it is obtained by finitely many applications of rules 1, 2 and 3.
- There are infinitely many n -variable Boolean formulas.

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n -variable Boolean Functions

- A Boolean function is a mapping that can be described by a Boolean formula.
- Given an n -variable Boolean formula F , the corresponding n -variable function $f : \mathbf{B}^n \rightarrow \mathbf{B}$ is defined as follows:
 1. If $F = b \in \mathbf{B}$, then the formula represents the **constant** function defined by
$$f(x_1, \dots, x_n) = b \quad \forall ([x_1], \dots, [x_n]) \in \mathbf{B}^n$$
 2. If $F = x_i$, then the formula represents the **projection** function defined by
$$f(x_1, \dots, x_n) = x_i \quad \forall ([x_1], \dots, [x_n]) \in \mathbf{B}^n$$
where $[x_k]$ denotes a valuation of variable x_k

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n -variable Boolean Functions

3. If the formula is of type either $G + H$, GH or G' , then the corresponding n -variable function is defined as follows
$$(g + h)(x_1, \dots, x_n) = g(x_1, \dots, x_n) + h(x_1, \dots, x_n)$$
$$(g \cdot h)(x_1, \dots, x_n) = g(x_1, \dots, x_n) \cdot h(x_1, \dots, x_n)$$
$$(g')(x_1, \dots, x_n) = g(x_1, \dots, x_n)'$$
for $\forall ([x_1], \dots, [x_n]) \in \mathbf{B}^n$
- The number of n -variable Boolean functions over a finite Boolean algebra $\mathbf{B}$ is *finite*.

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Example

- $\mathbf{B} = \{\underline{0}, \underline{1}, a, a'\}$
- Variable symbols:
 $\{x, y\}$
- 2-variable Boolean formula:
e.g., $a'x + ay'$
- 2-variable Boolean function: $f : \mathbf{B}^2 \rightarrow \mathbf{B}$
- Domain $\mathbf{B}^2 = \{(\underline{0}, \underline{0}), (\underline{0}, \underline{1}), \dots, (a, a)\}$

x	y	f
<u>0</u>	<u>0</u>	a
<u>0</u>	<u>1</u>	<u>0</u>
<u>0</u>	a'	a
<u>0</u>	a	<u>0</u>
<u>1</u>	<u>0</u>	<u>1</u>
<u>1</u>	<u>1</u>	a'
<u>1</u>	a'	<u>1</u>
<u>1</u>	a	a'
a	<u>0</u>	a
a	<u>1</u>	<u>0</u>
a	a'	a
a	a	<u>0</u>
a'	<u>0</u>	<u>1</u>
a'	<u>1</u>	a'
a'	a'	<u>1</u>
a'	a	a'

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Boole's Expansion Theorem

Theorem 1 If $f : \mathbf{B}^n \rightarrow \mathbf{B}$ is a Boolean function, then

$$f(x_1, \dots, x_n) = x'_1 f(\underline{0}, \dots, x_n) + x_1 f(\underline{1}, \dots, x_n)$$

for $\forall ([x_1], \dots, [x_n]) \in \mathbf{B}^n$

Proof. Case analysis of Boolean functions under the construction rules. Apply postulates of Boolean algebra.

- The theorem holds not only for two-element Boolean algebra (c.f. Shannon expansion)

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Minterm Canonical Form

Theorem 2 A function $f : \mathbf{B}^n \rightarrow \mathbf{B}$ is Boolean if and only if it can be expressed in the minterm canonical form

$$f(X) = \sum_{A \in \{0,1\}^n} f(A) \cdot X^A$$

where $X = (x_1, \dots, x_n) \in \mathbf{B}^n$, $A = (a_1, \dots, a_n) \in \{0,1\}^n$, and $X^A \equiv x_1^{a_1} \cdot x_2^{a_2} \dots x_n^{a_n}$ (with $x^0 \equiv x'$ and $x^1 \equiv x$)

Proof.

($\Rightarrow$) Follows from Boole's expansion theorem.

($\Leftarrow$) Examine the construction rules of Boolean functions.

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Example

f is **not** Boolean!

Proof. If f is Boolean, f can be expressed by $f(x) = x f(1) + x' f(0)$
 $= x + a x'$ from the minterm canonical form. However, substituting $x = a$ in the previous expression yields: $f(a) = a + a a'$
 $= a \neq 1$

x	f(x)
0	a
1	1
a'	a'
a	1

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Why Study General Boolean Algebra?

General algebras can't be avoided

$$f = x y + x z' + x' z$$

- Two-element view: $x, y, z \in \{0,1\}$ and $f \in \{0,1\}$
- General algebra view: f as a member of the Boolean algebra of 3-variable Boolean functions

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Why Study General Boolean Algebra?

General algebras are useful

- Two-element view: Truth tables include only 0 and 1.
- General algebra view: Truth tables contain any elements of $\mathbf{B}$.

J	K	Q	Q+
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
..	..	..	..

J	K	Q+
0	0	Q
0	1	0
1	0	1
1	1	Q'

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