

Introduction to Electronic Design Automation

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Model of Computation

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Model of Computation

- In system design, intended system behavior is translated into physical implementation
 - The physical implementation can be in hardware or software, in silicon or non-silicon (e.g., living cells)
 - How a system behaves or interacts with its environmental stimuli must be specified *formally*
- Model of computation (MoC) can be seen as the subject of devising/selecting effective “data structures” in describing system behaviors precisely and concisely
- MoC gives a **formal** way of describing system behaviors
 - It is useful in the **specification**, **synthesis** and **verification** of systems

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Model of Computation

- Outline
 - State transition systems
 - Finite automata / finite state machines
 - Real-time systems
 - Timed automata
 - Hybrid systems
 - Hybrid automata for hybrid systems, which exhibits both discrete and continuous dynamic behavior
 - Asynchronous systems
 - Petri nets for asynchronous handshaking
 - Signal processing systems
 - Dataflow process network for signal processing applications

(See Wikipedia for more detailed introduction)

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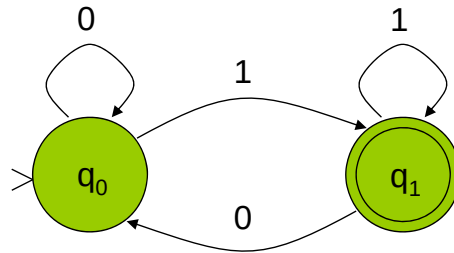
Modeling State Transition

□ Finite automata $A = (Q, q_0, F, \Sigma, \delta)$

- Q : states; q_0 : initial state; F : accepting states; Σ : input alphabet; $\delta: \Sigma \times Q \rightarrow Q$ transition
- Can be alternatively represented in *state diagram*

□ Finite automata are used as the recognizer of *regular language*

■ Example



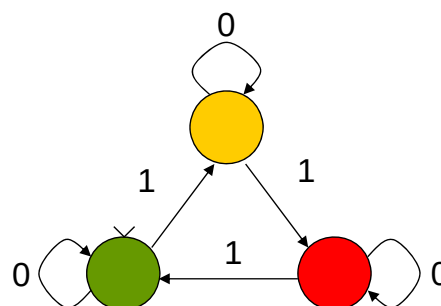
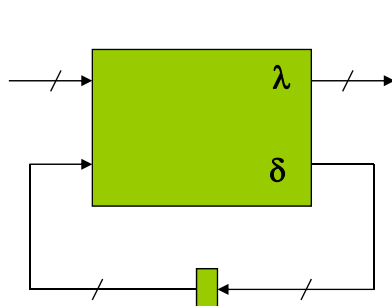
- The finite automaton accepts **all binary strings ended in a "1"**, i.e., which form the **language: $(0^*1^*)^*1$ or $\{0,1\}^*1$**

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Modeling State Transition (cont'd)

□ Finite state machine (FSM) $M = (Q, I, \Sigma, \Omega, \delta, \lambda)$

- Q : states; I : initial states; Σ : input alphabet; Ω : output alphabet; $\delta: \Sigma \times Q \rightarrow Q$ transition function; $\lambda: \Sigma \times Q \rightarrow \Omega$ (respectively $\lambda: Q \rightarrow \Omega$) output function for Mealy (respectively Moore) FSM
- Can be alternatively represented in *state transition graph* (STG) or *state transition table* (STG)
- E.g., vending machine, traffic light controller, elevator controller, Rubik's cube!, etc.



Σ	Q	Q'	Ω
0	G	G	go
0	Y	Y	yield
0	R	R	stop
1	G	Y	go
1	Y	R	yield
1	R	G	stop

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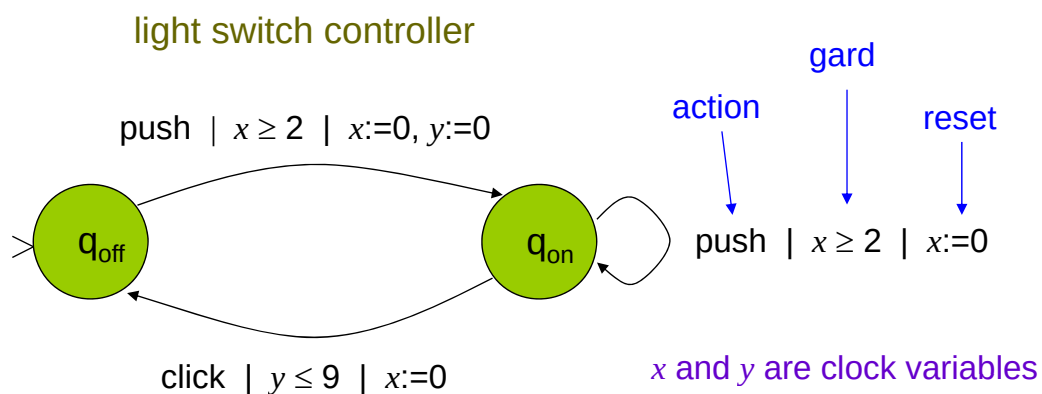
Modeling State Transition (cont'd)

- FSMs are often used as controllers in digital systems
 - E.g. data flow controller, ALU (arithmetic logic unit) controller, etc.
- Variants of FSM
 - Hierarchical FSM
 - Communicating FSM
 - ...

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Modeling Real-Time Systems

- Timed automata
 - Example



- Switch may be turned on whenever at least 2 time units has elapsed since last turn off
- Light switches off automatically after 9 time units

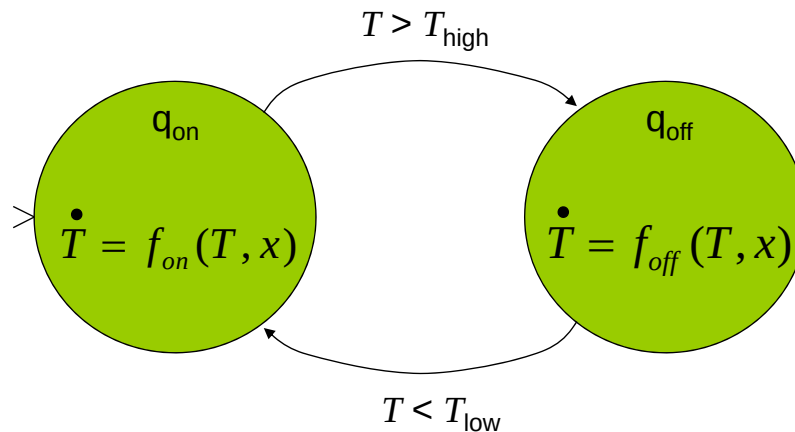
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Modeling Hybrid Systems

Hybrid automata

Example

temperature control system



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Modeling Asynchronous Systems

Petri net $P = (G, M_0)$

Petri net graph G is a bipartite weighted directed graph:

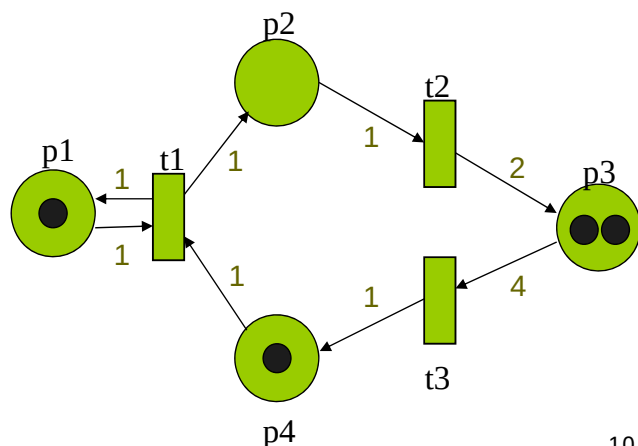
- Two types of nodes: *places* in circles and *transitions* in boxes

- Arcs: arrows labeled with weights indicating how many tokens are consumed or produced

- Tokens: black dots in places

Initial marking M_0

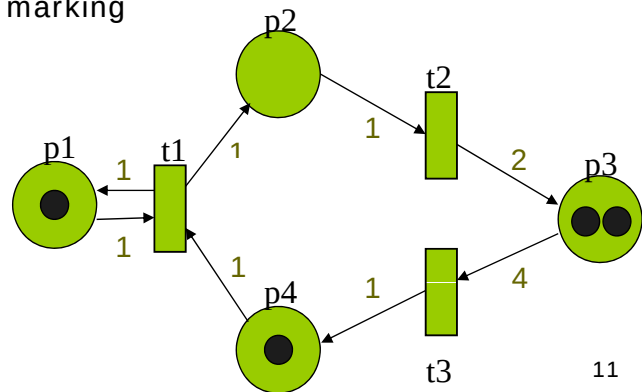
- Initial token positions



Modeling Asynchronous Systems (cont'd)

□ In a Petri net graph G ,

- **places**: represent distributed state by holding tokens
 - marking (state) M is an vector $(m_1, m_2, \dots, m_n)$, where m_i is the non-negative number of tokens in place p_i
 - initial marking M_0 is initial state
- **transitions**: represent actions/events
 - enabled transition: enough tokens in predecessors
 - firing transition: modifies marking

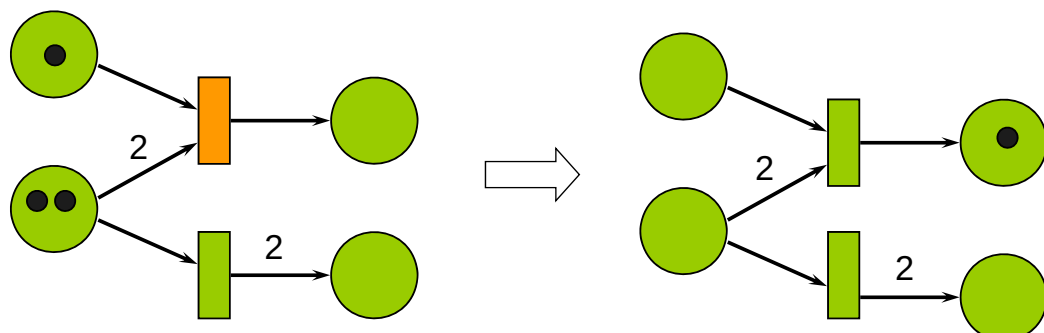


ref: EE249 lecture notes, UC Berkeley

Modeling Asynchronous Systems (cont'd)

□ A marking is changed according to the following rules:

- A transition is **enabled** if there are enough tokens in each input place
- An enabled transition **may or may not** fire (i.e. non-deterministic)
- The **firing** of a transition modifies marking by **consuming** tokens from the input places and **producing** tokens in the output places

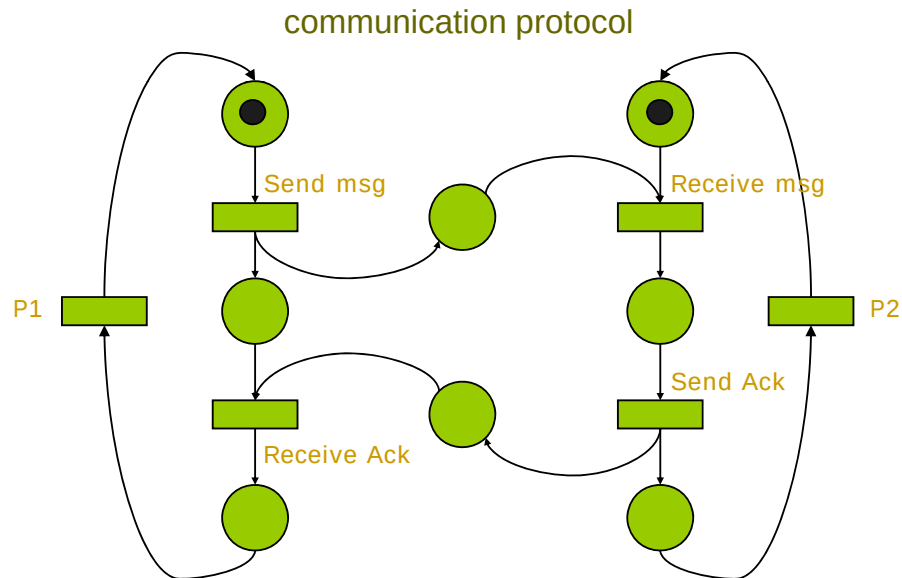


ref: EE249 lecture notes, UC Berkeley

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Modeling Asynchronous Systems (cont'd)

□ Example

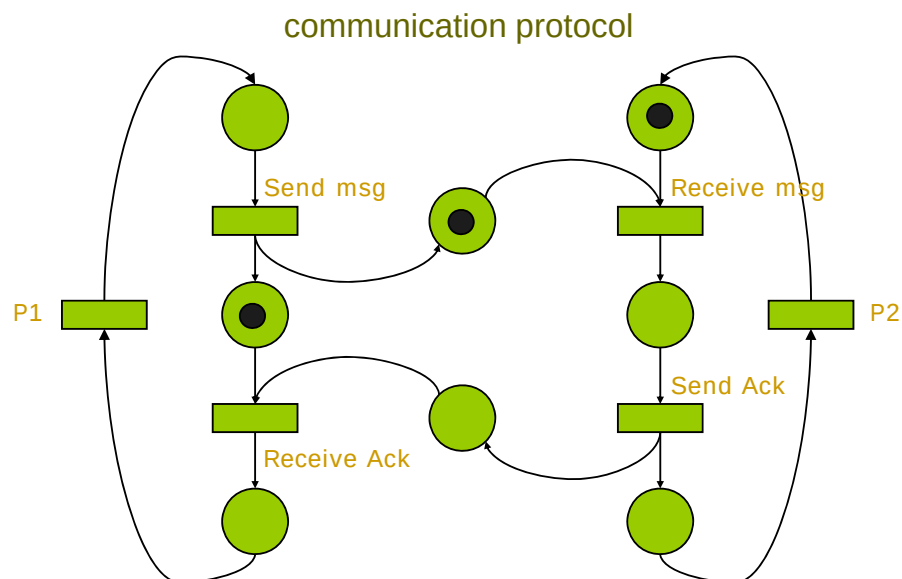


source: EE249 lecture notes, UC Berkeley

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Modeling Asynchronous Systems (cont'd)

□ Example

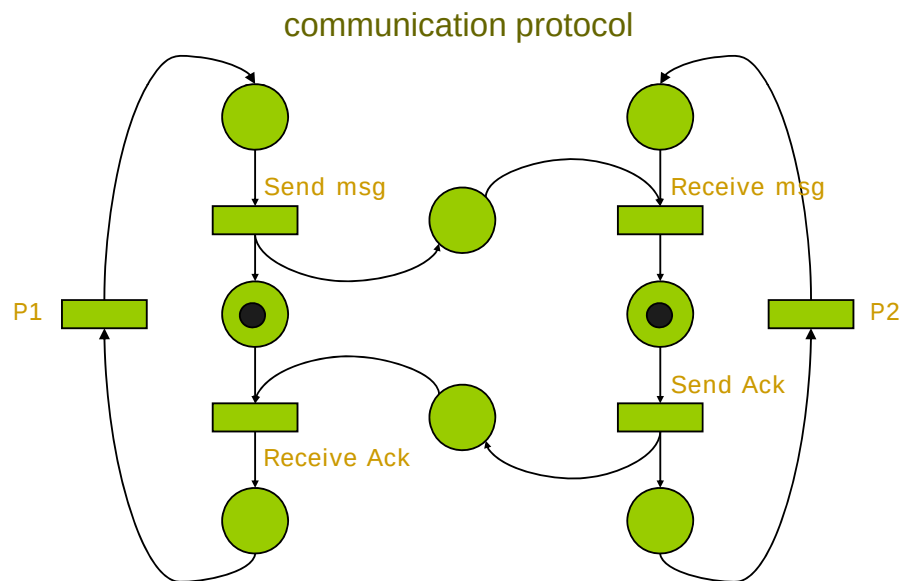


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Modeling Asynchronous Systems (cont'd)

□ Example

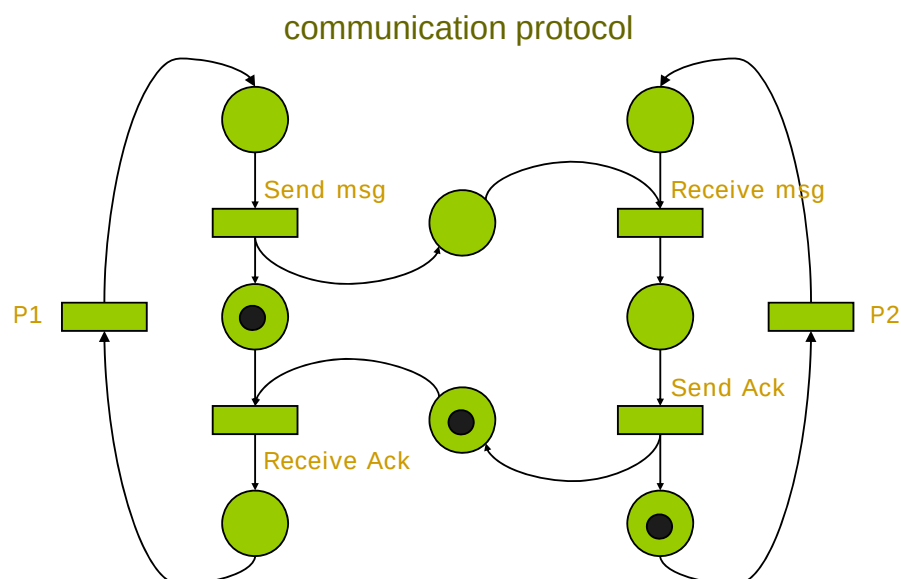


source: EE249 lecture notes, UC Berkeley

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Modeling Asynchronous Systems (cont'd)

□ Example

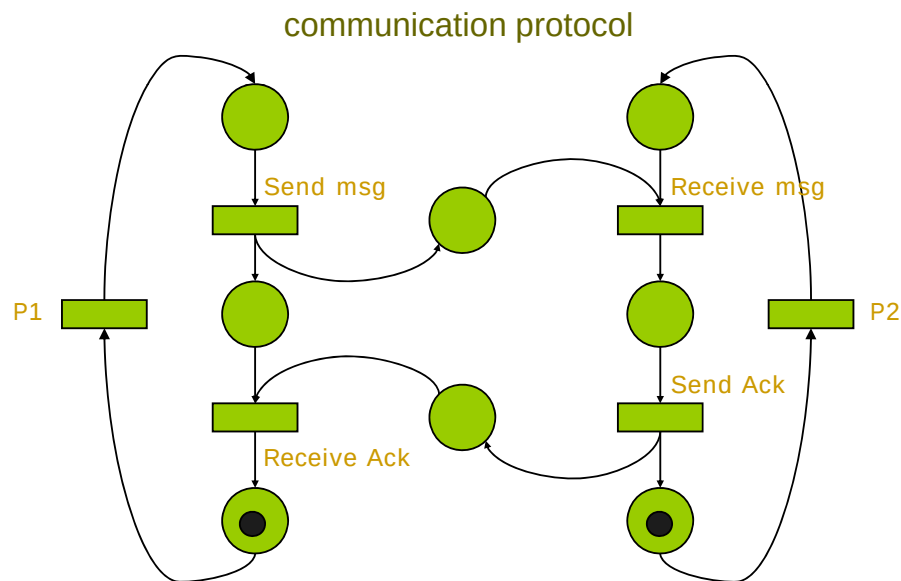


source: EE249 lecture notes, UC Berkeley

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Modeling Asynchronous Systems (cont'd)

□ Example

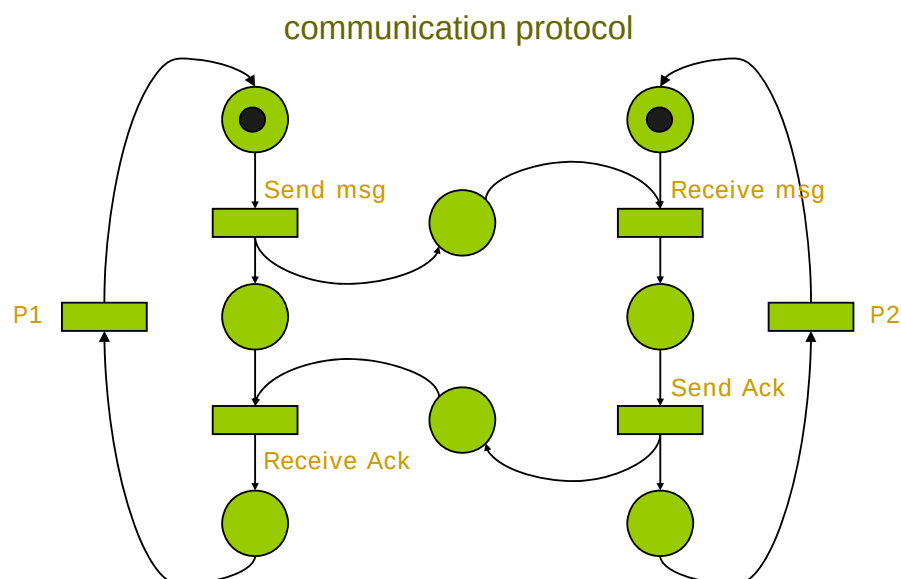


source: EE249 lecture notes, UC Berkeley

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Modeling Asynchronous Systems (cont'd)

□ Example



source: EE249 lecture notes, UC Berkeley

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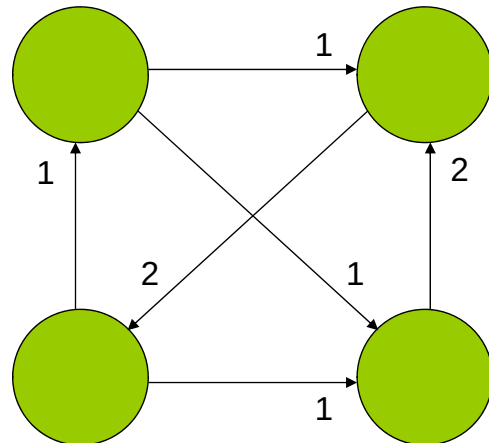
Modeling Signal Processing

- Data-flow process network

- Nodes represent **actors**; arcs represent **FIFO queues**

- **Firing rules** are specified on arcs

- Actors respect firing rules that specify how many tokens must be available on every input for an actor to fire. When an actor fires, it consumes a finite number of tokens and produces also a finite number of output tokens.



ref: <http://www.create.ucsb.edu/~xavier/Thesis/html/node38.html>

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MoC in System Construction

- There are many other models of computation tailored for specific applications

- Can you devise a new computation model in some domain?

- Hierarchical modeling combined with several different models of computation is often necessary

- By using a proper MoC, a system can be specified formally, and further synthesized and verified

- In the sequel of this course, we will be focusing on FSMs mainly

High Level Synthesis

High-level synthesis



Logic synthesis



Physical design

Slides are by Courtesy of Prof. Y.-W. Chang

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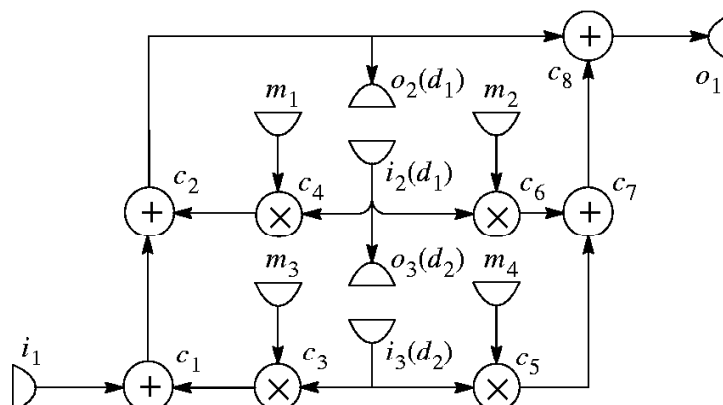
High Level Synthesis

□ Course contents

- Hardware modeling
- Data flow
- Scheduling/allocation/assignment

□ Reading

- Chapter 5



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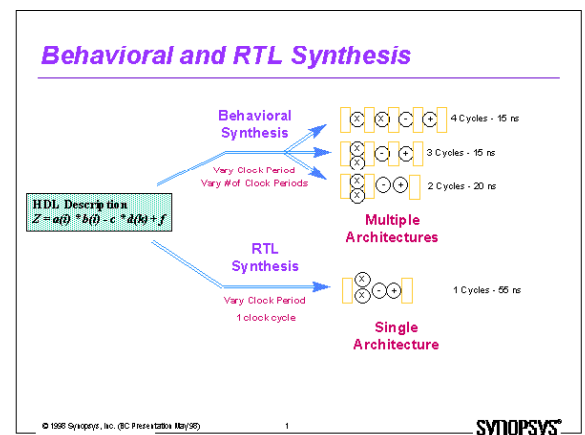
High Level Synthesis

- ❑ Hardware-description language (HDL) synthesis
 - Starts from a **register-transfer level** (RTL) description; **circuit behavior in each clock cycle is fixed**
 - Uses logic synthesis techniques to optimize the design
 - Generates a netlist
- ❑ High-level synthesis (HLS), also called architectural or behavioral synthesis
 - Starts from an abstract behavioral description
 - Generates an RTL description
 - It normally has to perform the trade-off between the **number of cycles** and the **hardware resources** to fulfill a task

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HL Synthesis vs. RTL Synthesis

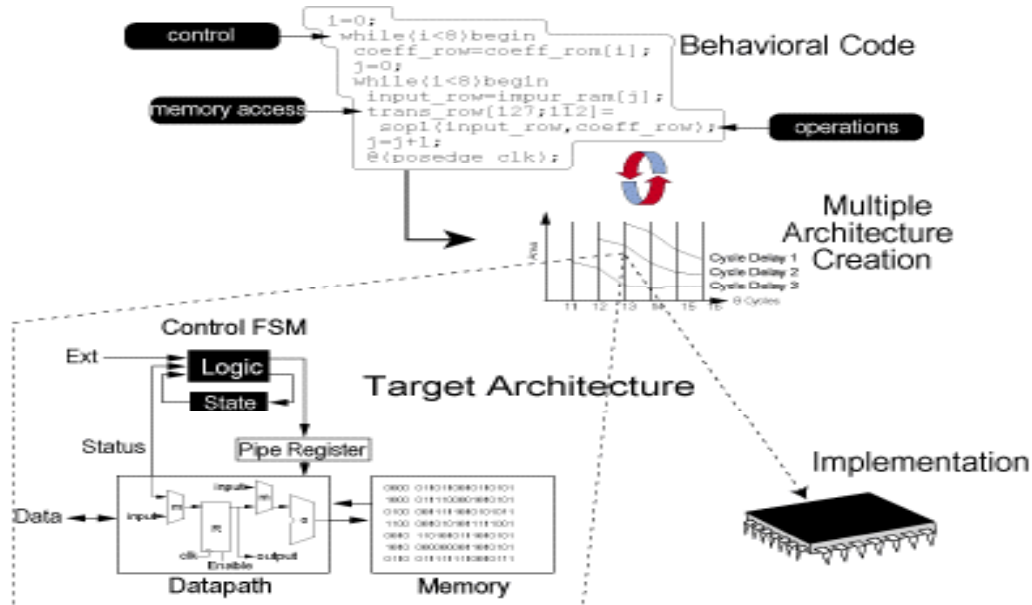
- ❑ RTL synthesis implements all functionality within a single clock cycle
- ❑ HL synthesis automatically allocates the functionality across multiple clock cycles



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Output of High Level Synthesis

- Behavioral Compiler creates a design that consists of a datapath, memory I/O and a control FSM



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Benefits of High Level Synthesis (1)

- Quick specification and verification
 - Specify behavioral HDL easily, since it's intuitive and natural to write
 - Save time -- behavioral HDL code is up to 10 times shorter than equivalent RTL
 - Simulate orders of magnitude faster because of the higher level of abstraction
 - Reuse designs more readily by starting with a more abstract description
- Reduce design time
 - Model hardware and software components of system concurrently
 - Easily implement algorithms in behavioral HDL and generate RTL code with a behavioral compiler
 - Verify hardware in system context at various levels of abstraction

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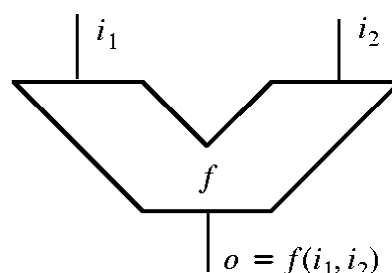
Benefits of High Level Synthesis (2)

- Explore architectural trade-offs
 - Create multiple architectures from a single specification
 - Trade-off throughput and latency using high-level constraints
 - Analyze various combinations of technology-specific datapath and memory resources
 - Evaluate cost/performance of various implementations rapidly
- Automatically infer memory and generate FSM
 - Specify memory reads and writes
 - Schedule memory I/O, resolve conflicts by building control FSM
 - Trade-off single-ported (separate registers) vs. multi-ported memories (register files)
 - Generate a new FSM

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Hardware Models for HL Synthesis

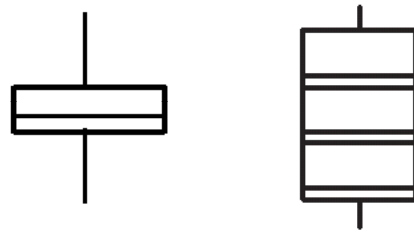
- All HLS systems need to restrict the target hardware
 - Otherwise search space is too large
- All synthesis systems have their own peculiarities, but most systems generate **synchronous** hardware and build it with **functional units**:
 - A functional unit can perform one or more computations, e.g., addition, multiplication, comparison, ALU



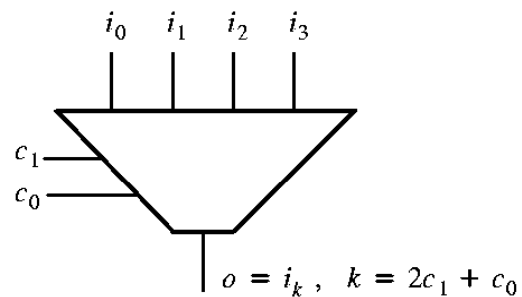
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Hardware Models

□ **Registers:** they store inputs, intermediate results and outputs; sometimes several registers are taken together to form a register file



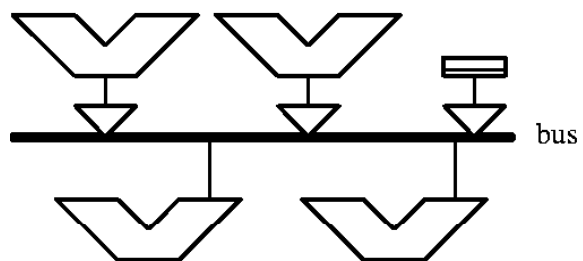
□ **Multiplexers:** from several inputs, one is passed to the output



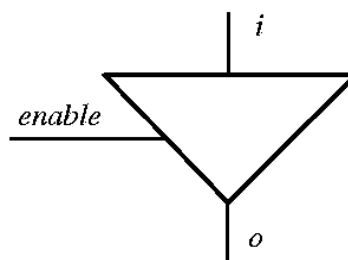
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Hardware Models (cont'd)

□ **Buses:** a connection shared between several hardware elements, such that only one element can write data at a specific time



□ **Tri-state drivers:** control the exclusive writing on the bus



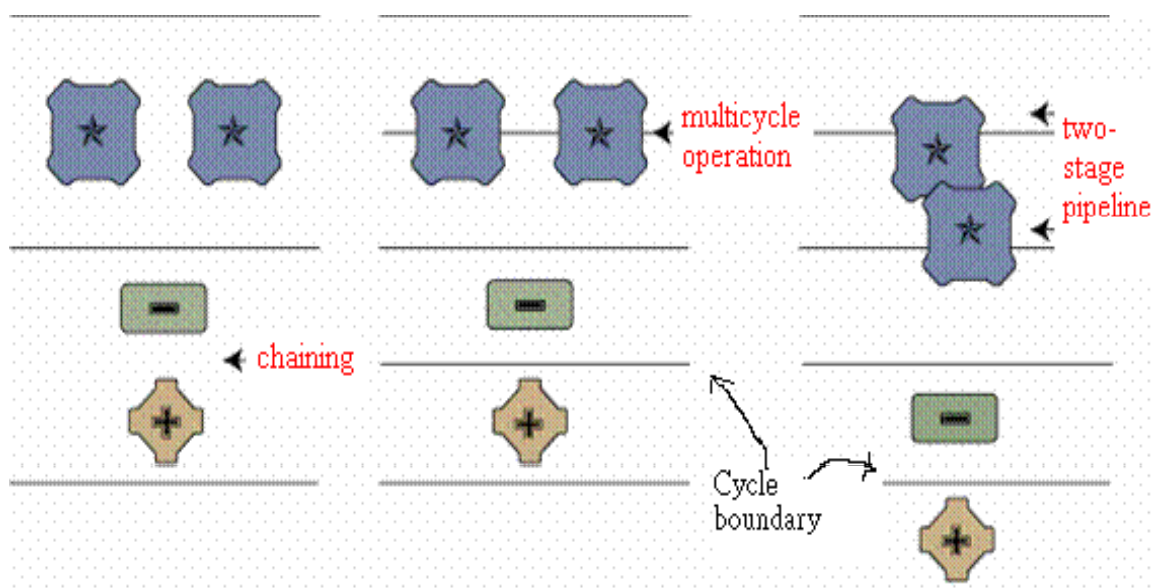
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Hardware Models (cont'd)

- ❑ Parameters defining the hardware model for the synthesis problem:
 - **Clocking strategy:** e.g. single or multiple phase clocks
 - **Interconnect:** e.g. allowing or disallowing buses
 - **Clocking of functional units:** allowing or disallowing
 - ❑ multicycle operations
 - ❑ operation chaining (multiple operations in one cycle)
 - ❑ pipelined units

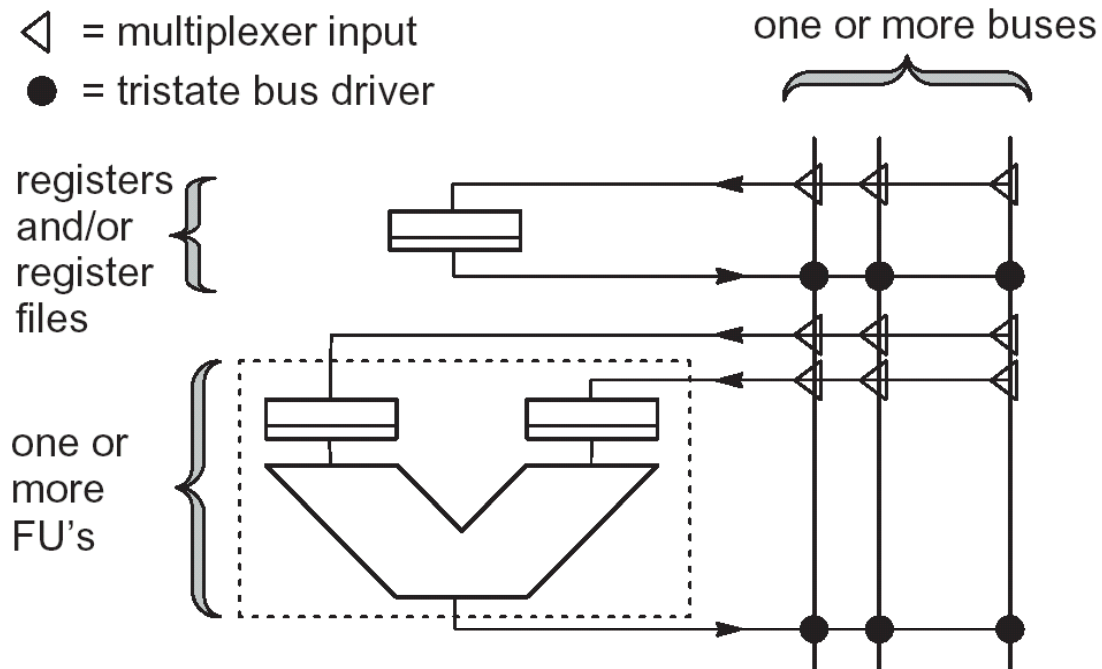
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Chaining, Multicycle Operation, Pipelining



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Example of a HLS Hardware Model



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Hardware Concepts: Data Path + Control

- Hardware is normally partitioned into two parts:
 - **Data path:** a network of functional units, registers, multiplexers and buses
 - The actual “computation” takes place in the data path
 - **Control:** the part of the hardware that takes care of having the data present at the right place at a specific time, of presenting the right instructions to a programmable unit, etc.
- High-level synthesis often concentrates on data-path synthesis
 - The control part is then realized as a finite state machine or in microcode

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Steps of High Level Synthesis

- ❑ **Preprocess** the design with high-level optimization
 - Code motion
 - Common subexpression elimination
 - Loop unrolling
 - Constant propagation
 - Modifications taking advantage of associativity and distributivity, etc.
- ❑ **Transform** the optimized design into intermediate format (internal representation) which reveals more structural characteristics of the design
- ❑ **Optimize** the intermediate format
 - Tree height reduction
 - Behavior retiming
- ❑ **Allocate** the required resources to implement the design
 - Also called module selection
- ❑ **Schedule** each operation to be performed at certain time such that no precedence constraint is violated
- ❑ **Assign (bind)** each operation to a specific functional unit and each variable to a register

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HLS Optimization Criteria

- ❑ Typically, in terms of speed, area, and power consumption
- ❑ Optimization is often constrained
 - Optimize area when the minimum speed is given $\Rightarrow$ time-constrained synthesis
 - Optimize speed when a maximum for each resource type is given $\Rightarrow$ resource-constrained synthesis
 - ❑ E.g. power-constrained synthesis
 - Minimize power dissipation for a given speed and area requirement $\Rightarrow$ time- and area- constrained synthesis

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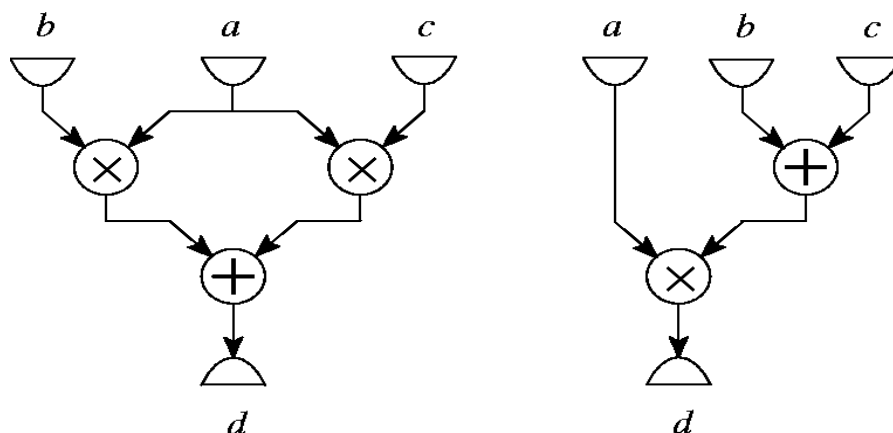
Input Format

- ❑ The algorithm, which is the input to a high-level synthesis system, is often provided in textual form either
 - in a conventional programming language, such as C, C++, SystemC, or
 - in a hardware description language (HDL), which is more suitable to express the parallelism present in hardware.
- ❑ The description has to be parsed and transformed into an internal representation and thus conventional compiler techniques can be used.

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Example of HL Optimization

- ❑ Applying the distributive law to reduce resource requirement



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Internal Representation

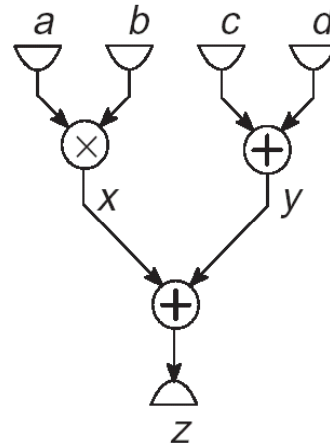
- Most systems use some form of a **data-flow graph (DFG)**

- A DFG may or may not contain information on control flow

- A data-flow graph is built from

- vertices (nodes): representing computation, and
- edges: representing **precedence** relations

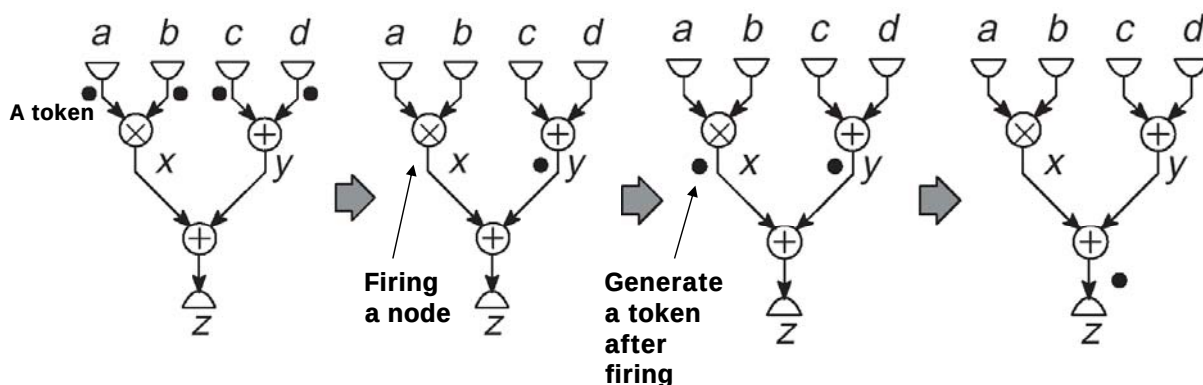
```
x := a * b; y := c + d;  
z := x + y;
```



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Token Flow in a DFG

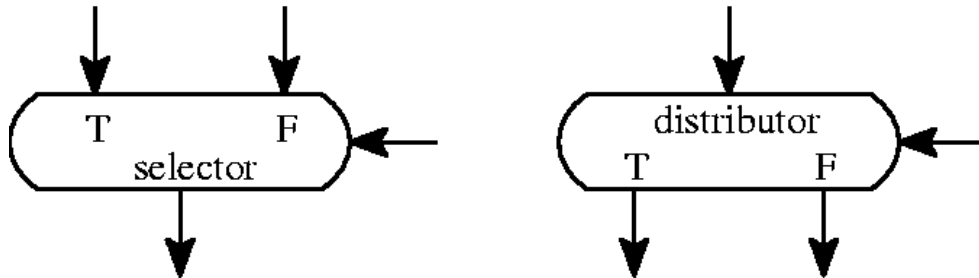
- A node in a DFG **fires** when all **tokens** are present at its inputs
- The input tokens are consumed and an output token is produced (like in Petri nets)



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Conditional Data Flow

- Conditional data flow by means of two special nodes:



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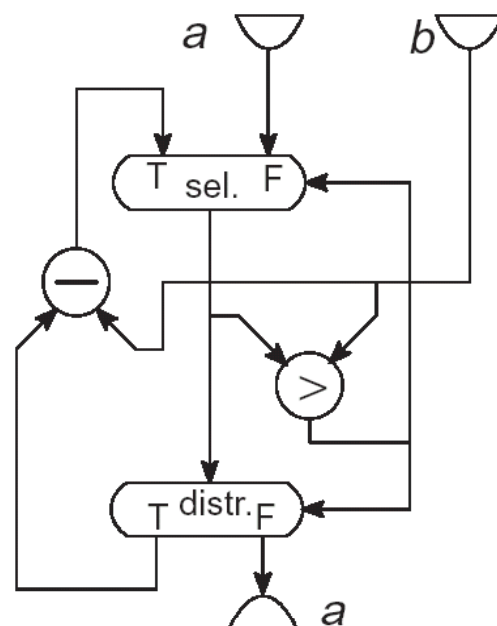
Explicit Iterative Data Flow

- Selector and distributor nodes can be used to describe iteration

- Example

```
while (a > b)
  a ← a - b;
```

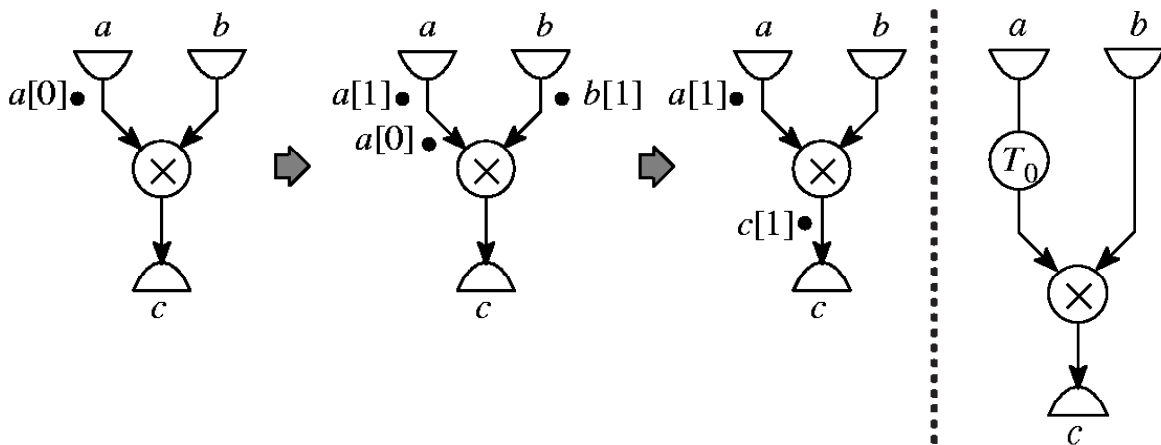
- Loops require careful placement of initial tokens on edges



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Implicit Iterative Data Flow

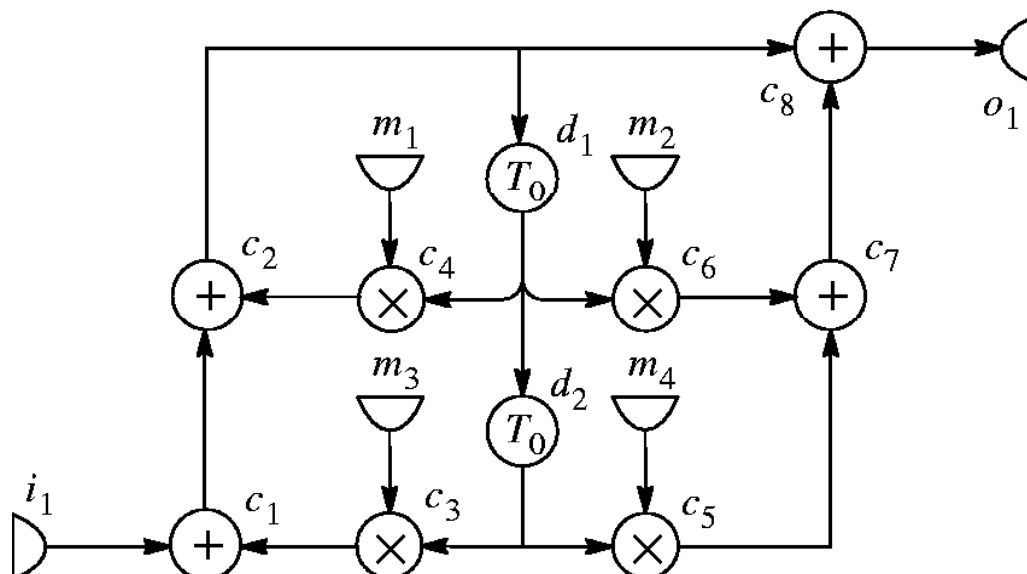
- Iteration implied by regular input stream of tokens
- Initial tokens act as buffers
- Delay elements instead of initial tokens



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Iterative DFG Example

a second-order filter section



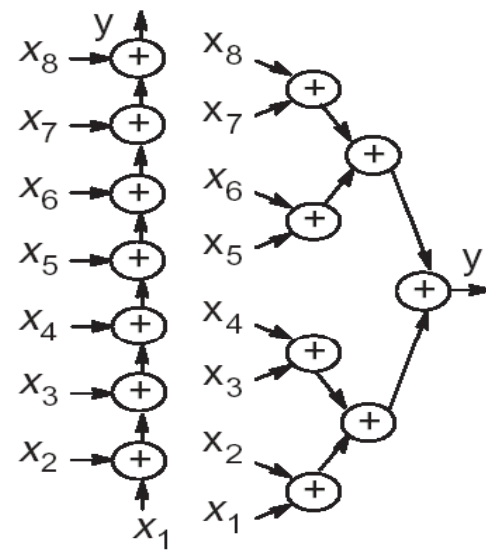
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Optimization of Internal Representation

- Restructuring data and control flow graphs prior to the actual mapping onto hardware

- Examples:

- Replacing chain of adders by a tree
- Behavior retiming

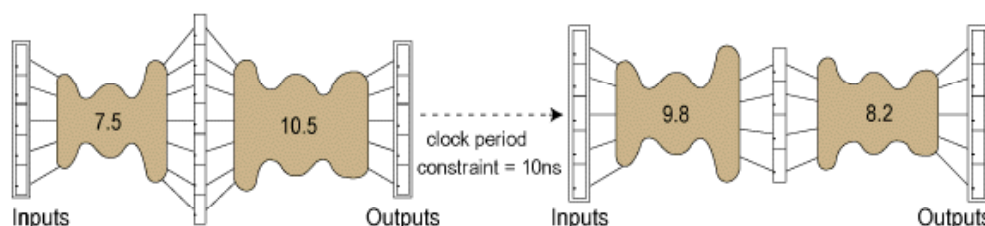
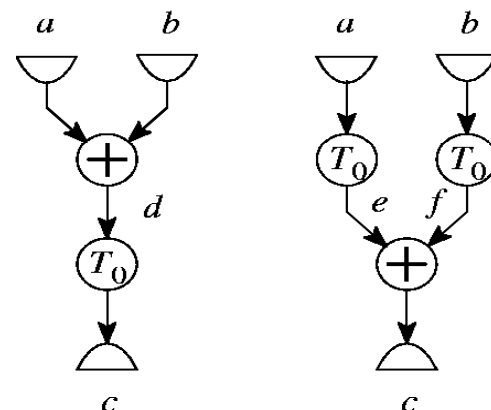


Tree height reduction

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Behavior Retiming (BRT)

- By moving registers through logic and hierarchical boundaries, BRT reduces the clock period with minimum area impact



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Effectiveness of Behavior Retiming

Synopsys exp:

DESIGN TYPE	RTL DESIGN		BEHAVIORAL RETIMING		SUMMARY
	SPEED (CLOCK PERIOD)	GATES	SPEED (CLOCK PERIOD)	GATES	
Control	44 ns	10,913 gates	30.6 ns	11,313 gates	30% faster, 4% more area
Control	23.1 ns	3,598 gates	19.6 ns	4,575 gates	15% faster, 27% more area
Control	28.6 ns	3,585 gates	28.6 ns	3,369 gates	same speed, 6% less area
Dataflow & Control	17 ns	28,900 gates	12.5 ns	30,100 gates	26% faster, 4% more area
Dataflow & Control	16 ns	7,620 gates	13 ns	8,019 gates	20% faster, 5% more area
Dataflow	22 ns	4,990 gates	18.5 ns	5,109 gates	16% faster, 2% more area
Dataflow	28 ns	31,226 gates	26 ns	32,032 gates	8% faster, 2% more area
Dataflow	26.2 ns	14,351 gates	23.6 ns	13,847 gates	10% faster, 4% less area
Dataflow	25.9 ns	16,798 gates	20.8 ns	15,550 gates	20% faster, 7% less area
Dataflow	45 ns	28,705 gates	26 ns	30,987 gates	42% faster, 8% more area

- ❑ RTL designs have a single clock net and were synthesized into gates using Synopsys Design Compiler
- ❑ Design type: dataflow implies significant number of operators; control implies state machine dominated

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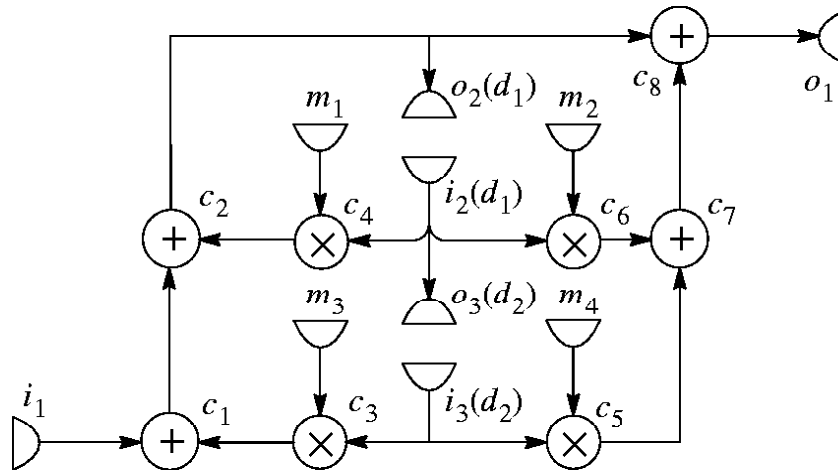
HLS Subtasks: Allocation, Scheduling, Assignment

- ❑ Subtasks in high-level synthesis
 - **Allocation (Module selection):** specify the hardware resources that will be necessary
 - **Scheduling:** determine for each operation the time at which it should be performed such that no precedence constraint is violated
 - **Assignment (Binding):** map each operation to a specific functional unit and each variable to a register
- ❑ Remarks:
 - Though the subproblems are strongly interrelated, they are often solved separately. However, to attain a better solution, an iterative process executing these three subtasks must be performed.
 - Most scheduling problems are NP-complete $\Rightarrow$ heuristics are used

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Example of High Level Synthesis

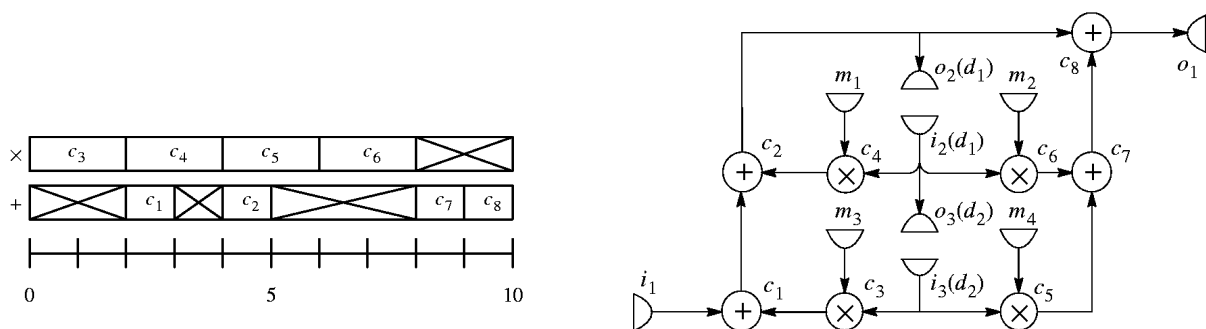
- The following couple of slides shows an example of scheduling and binding of a design with given resource allocation
- Given the second-order filter which is first made acyclic:



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Example of Scheduling

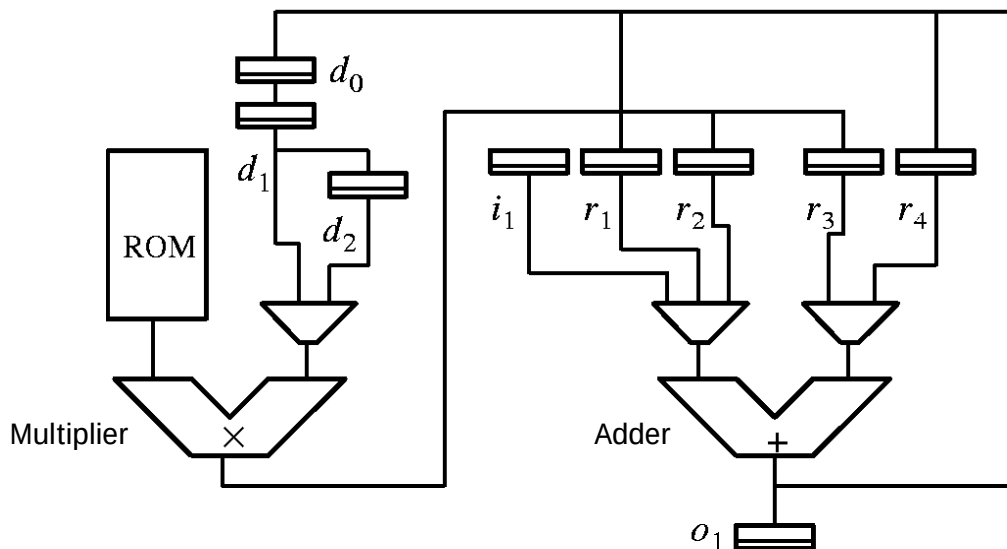
- The schedule and operation assignment with an allocation of one adder and one multiplier:



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Binding for Data Path Generation

- The resulting data path after register assignment
 - The specification of a controller would complete the design



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Resource Allocation Problem

- This problem is relatively simple. It simply decides the kinds of hardware resources (hardware implementation for certain functional units such as adder, multiplier, etc.) and the quantity of these resources.
 - For example two adders, one multiplier, 4 32-bit registers, etc. for a certain application
- The decision made in this step has a profound influence on the scheduling which under the given resource constraints decides the time when an operation should be executed by a functional unit
- This step set an upper bound on the attainable performance.

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Problem Formulation of Scheduling

- Input consists of a DFG $G(V, E)$ and a library $\mathcal{R}$ of resource types
- There is a fixed mapping from each $v \in V$ to some $r \in \mathcal{R}$; the execution delay $\delta(v)$ for each operation is therefore known
- The problem is time-constrained; the available execution times are in the set

$$\mathcal{T} = \{0, 1, \dots, T_0 - 1\}.$$

- A schedule $\sigma: V \rightarrow T$ maps each operation to its starting time; for each edge $(v_i, v_j) \in E$, a schedule should respect: $\sigma(v_j) \geq \sigma(v_i) + \delta(v_i)$.
- Given the resource type cost $\omega(r)$ and the requirement function $N_r(\sigma)$, the cost of a schedule σ is given by:

$$\sum_{r \in \mathcal{R}} \omega(r) N_r(\sigma).$$

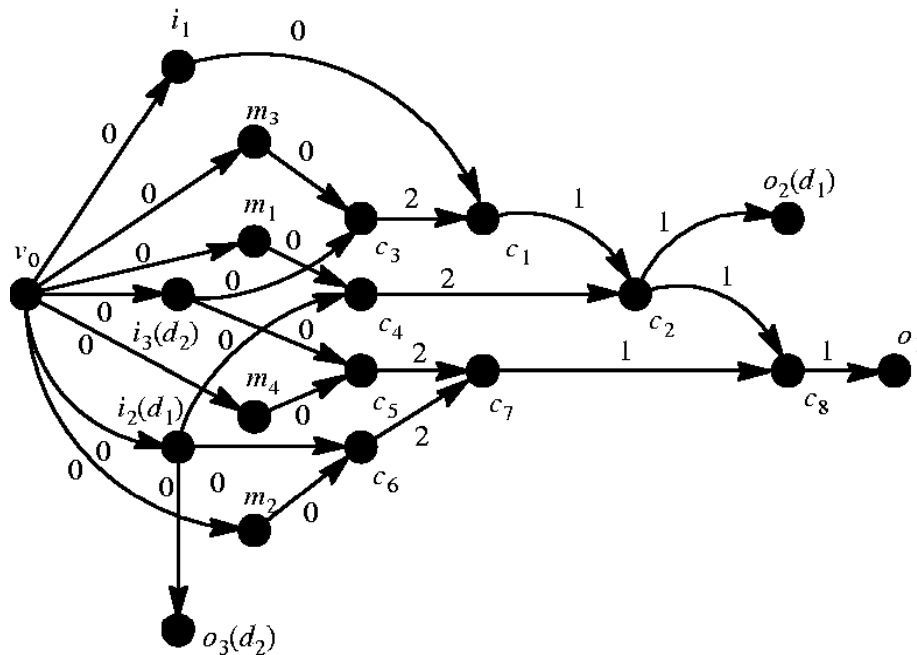
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ASAP Scheduling

- **As soon as possible (ASAP) scheduling** maps an operation to the earliest possible starting time not violating the precedence constraints
- Properties:
 - It is easy to compute by finding the longest paths in a directed acyclic graph
 - It does not make any attempt to optimize the resource cost

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Graph for ASAP Scheduling



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Mobility Based Scheduling

- Compute both the ASAP and **ALAP** (as late as possible) schedules σ_S and σ_L
- For each $v \in V$, determine the scheduling range $[\sigma_S(v), \sigma_L(v)]$
- $\sigma_L(v) - \sigma_S(v)$ is called the **mobility** of v
- Mobility-based scheduling tries to find the best position within its scheduling range for each operation

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Simple Mobility Based Scheduling

- A partial schedule $\tilde{\sigma} : V \rightarrow [\mathcal{T}, \mathcal{T}]$ assigns a scheduling range to each $v \in V$,

$$\tilde{\sigma}(v) = [\tilde{\sigma}_{\min}(v), \tilde{\sigma}_{\max}(v)]$$

- Finding a schedule can be seen as the generation of a sequence of partial schedules $\tilde{\sigma}^{(0)}, \dots, \tilde{\sigma}^{(n)}$.

“determine $\tilde{\sigma}^{(0)}$ by computing σ_S and σ_L ”;

$k \leftarrow 0$;

while (“there are unscheduled operations”) {

$v \leftarrow$ “one of the nodes with lowest mobility”;

 “schedule v at some time that optimizes the current resource utilization”;

 “determine $\tilde{\sigma}^{(k+1)}$ by updating the scheduling ranges of the unscheduled nodes”;

$k \leftarrow k + 1$

}

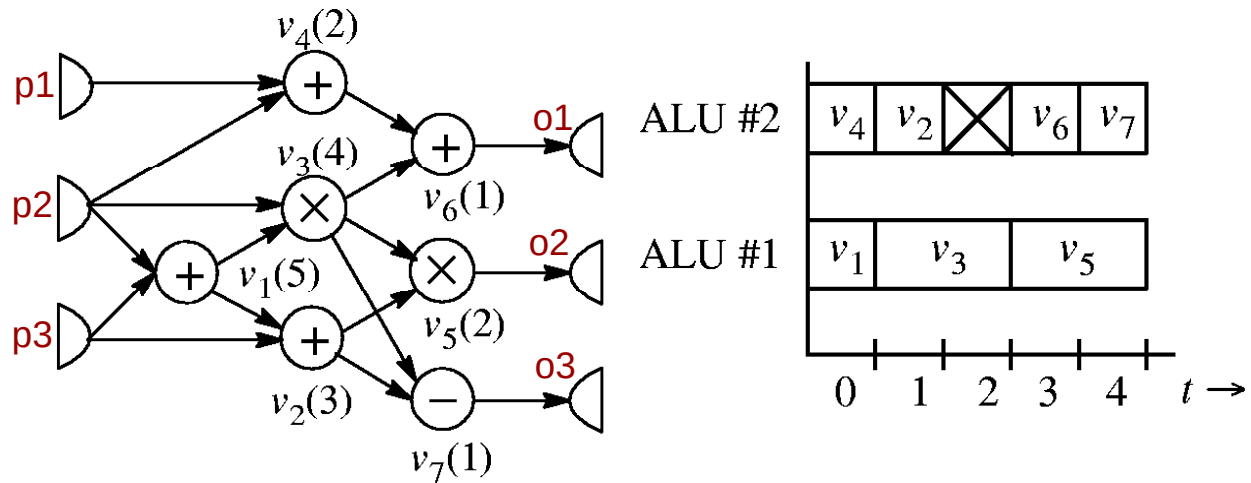
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List Scheduling

- A resource-constrained scheduling method
- Start at time zero and increase time until all operations have been scheduled
 - Consider the precedence constraint
- The ready list L_t contains all operations that can start their execution at time t or later
- If more operations are ready than there are resources available, use some priority function to choose, e.g. the longest-path to the output node
 $\Rightarrow$ critical-path list scheduling

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List Scheduling Example



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Assignment Problem

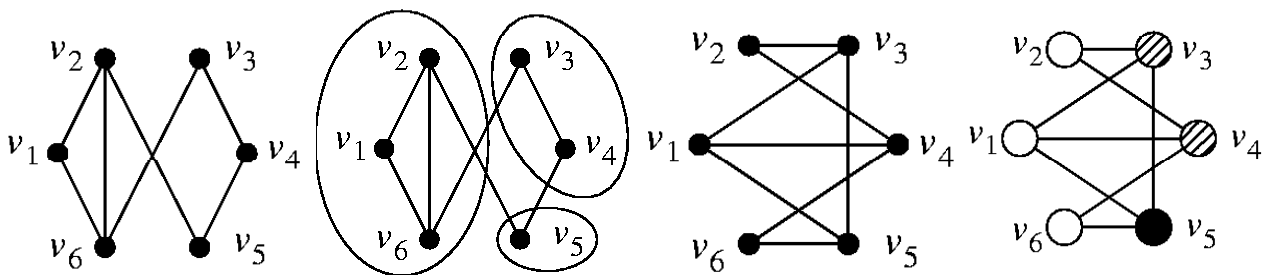
- Subtasks in assignment:
 - operation-to-FU assignment
 - value grouping
 - value-to-register assignment
 - transfer-to-wire assignment
 - wire to FU-port assignment
- In general: task-to-agent assignment

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Compatibility and Conflict Graphs

□ Clique partitioning gives an assignment in a **compatibility graph**

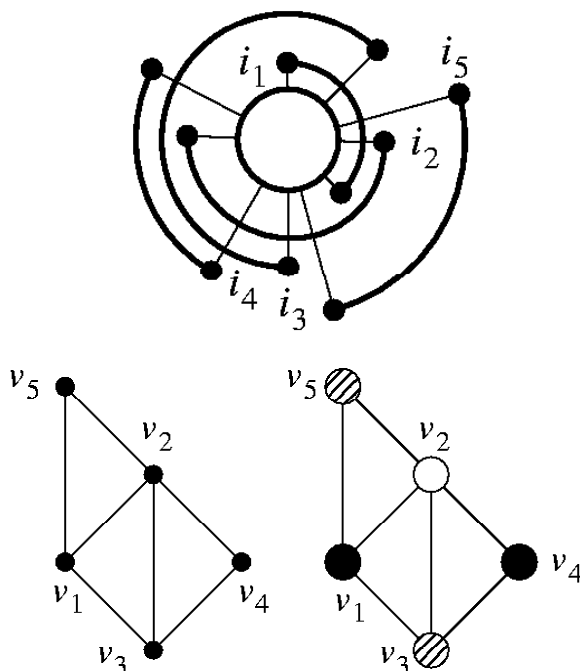
□ Graph coloring gives an assignment in the complementary **conflict graph**



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Assignment Problem

- Assumption: assignment follows scheduling.
- The claim of a task on an agent is an interval $\Rightarrow$ minimum resource utilization can be found by **left-edge** algorithm.
- In case of iterative algorithm, interval graph becomes **circular-arc** graph $\Rightarrow$ optimization is NP-complete.



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Tseng and Sieworek's Algorithm

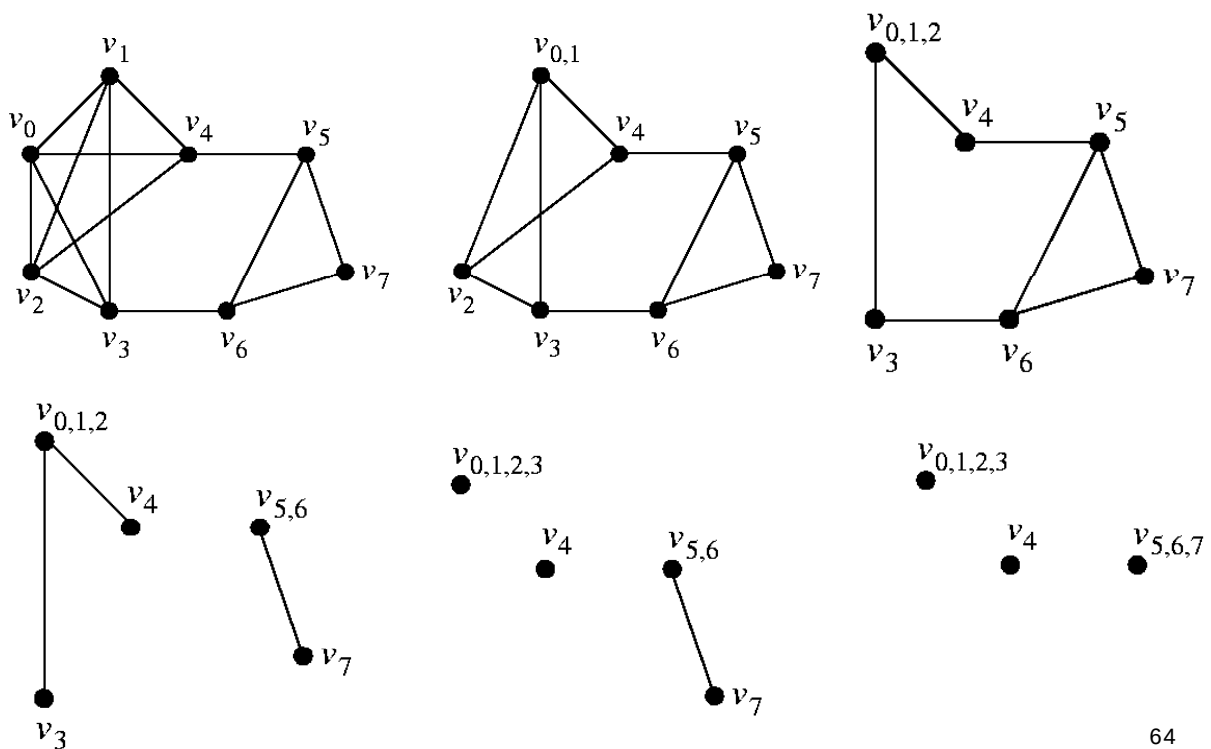
```

 $k \leftarrow 0;$ 
 $G_c^k(V_c^k, E_c^k) \leftarrow G_c(V_c, E_c);$ 
while ( $E_c^k \neq \emptyset$ ) {
    "find  $(v_i, v_j) \in E_c^k$  with largest set of common neighbors";
     $N \leftarrow$  "set of common neighbors of  $v_i$  and  $v_j$ ";
     $s \leftarrow i \cup j;$ 
     $V_c^{k+1} \leftarrow V_c^k \cup \{v_s\} \setminus \{v_i, v_j\};$ 
     $E_c^{k+1} \leftarrow \emptyset;$ 
    for each  $(v_m, v_n) \in E_c^k$ 
        if ( $v_m \neq v_i \wedge v_m \neq v_j \wedge v_n \neq v_i \wedge v_n \neq v_j$ )
             $E_c^{k+1} \leftarrow E_c^{k+1} \cup \{(v_m, v_n)\};$ 
    for each  $v_n \in N$ 
         $E_c^{k+1} \leftarrow E_c^{k+1} \cup \{(v_n, v_s)\};$ 
     $k \leftarrow k + 1;$ 
}

```

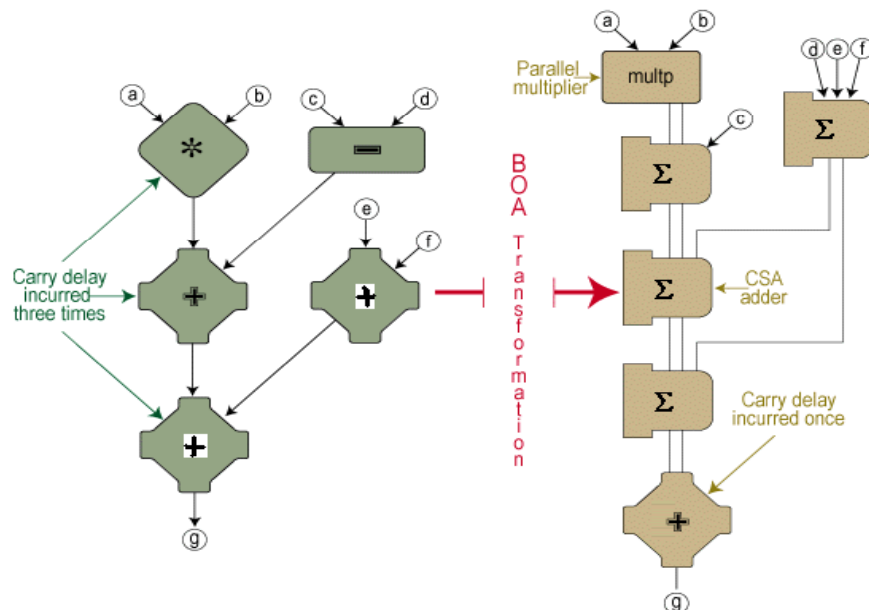
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Clique-Partitioning Example



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Example of Behavior Optimization



Behavior Optimization of Arithmetic Circuit (BOA)

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Effectiveness of BOA

Synopsys example

DESIGN TYPE	RTL DESIGN	BOA	SUMMARY
Motion estimation	23.6 ns 12,793 gates	20.2 ns 12,215 gates	14% faster, 5% less area
Graphics interpolation	19.2 ns 3,507 gates	17.5 ns 2,952 gates	9% faster, 16% less area
Color space conversion and scaling	16 ns 35,866 gates (manual CSA implementation)	14.9 ns 34,397 gates	7% faster, 4% less area
Sum of 9 operands	7.7 ns 1,418 gates	5.3 ns 1,307 gates	31% faster, 8% less area
$a * b + 1$	11.6 ns 2,577 gates	9.3 ns 2,524 gates	20% faster, 2% less area
$a * 4104$ (0100000100000100)	4.4 ns 759 gates	3.1 ns 449 gates	30% faster, 40% less area
$a * 3E3E$ (0011111000111110)	5.7 ns 927 gates	4.6 ns 709 gates	19% faster, 23% less area
$a * b + c$	11.0 ns 2,707 gates	10.0 ns 2,689 gates	9% faster, same area
$a * b + c * d + e * f$	14.2 ns 7,435 gates	12.8 ns 7,110 gates	10% faster, 4% less area
Sum of 16 operands	8.1 ns 2,836 gates	6.7 ns 2,123 gates	17% faster, 25% less area

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